

**MATH 61-02: WORKSHEET 2 (§1.4-1.5)**

- (W1) In class, you have been learning the building blocks of logic. In this problem, we explore some of the *replacement rules* of propositional logic. These are steps that allow us to replace a logical sentence with an equivalent one, knowing that the truth value remains intact in the course of the replacement.
- (a) One replacement rule says that we can replace  $P \wedge (Q \vee R)$  with  $(P \wedge Q) \vee (P \wedge R)$ . Verify that these two statements are equivalent using a truth table.

- (b) Another replacement rule says that  $P \Leftrightarrow Q$  can be replaced by  $(P \wedge Q) \vee (\sim P \wedge \sim Q)$ . Verify that this is true through a truth table, and explain in words why this would be true.

(W2) Calculus is about limits. However, you've probably never seen a formal definition. In this problem, we look at a mathematically quantified definition of the limit of a sequence.

- (a) Let  $(x_n) = x_1, x_2, \dots$  be a sequence of real numbers. We call  $x$  the *limit* of the sequence  $(x_n)$  if the following condition holds:

$$\forall \epsilon > 0, \exists N \in \mathbb{N} \text{ s.t. } \forall n > N, |x_n - x| < \epsilon.$$

Parse this statement in English and give some kind of intuitive explanation as to what this definition means.

- (b) Using quantifiers, give the negation of the above statement.

- (c) This definition of a limit can be expanded to define the limit of a function: Let  $f : \mathbb{R} \rightarrow \mathbb{R}$  be a function. Then the precise mathematical meaning of  $\lim_{x \rightarrow c} f(x) = L$  is

$$\forall \epsilon > 0, \exists \delta > 0 \text{ s.t. } |x - c| < \delta \implies |f(x) - L| < \epsilon.$$

Can you write down a similar expression for  $\lim_{x \rightarrow \infty} f(x) = L$ ?

- (d) Use quantifiers to express Fermat's Last Theorem:  
*No integers  $a, b, c$  satisfy the equation  $a^n + b^n = c^n$  for any integer exponent  $n$  greater than 2.*

(W3) In class, you saw the Sheffer stroke, a logical operator ("NAND") defined as follows:

<b>P</b>	<b>Q</b>	<b><math>P \uparrow Q</math></b>
T	T	F
T	F	T
F	T	T
F	F	T

Recall that  $P \uparrow P = \sim P$ . Work out expressions for both  $P \wedge Q$  and  $P \vee Q$  using only  $P, Q, \uparrow$ , and parentheses.