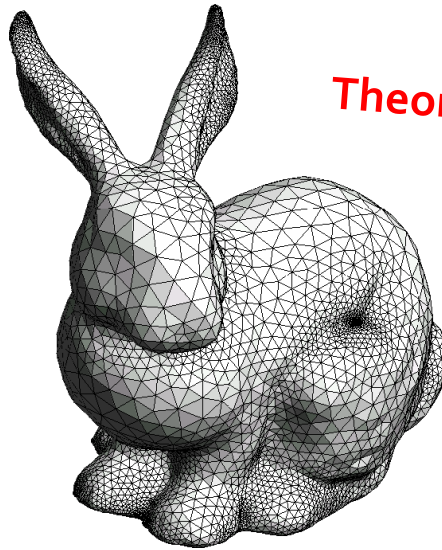


Tutorial on **Optimal Transport**

Justin Solomon

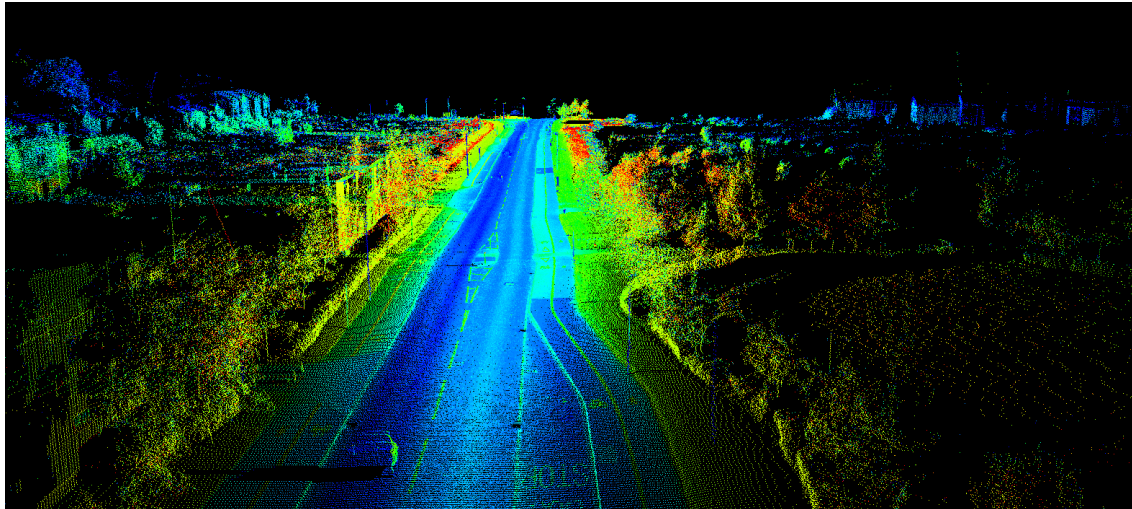


Motivation



Theory

Practice



Potential VRDI Motivation

- Reasonable notion of compactness:
Travel time to voting booth
- Connection to **Voronoi** and **power diagrams**
- Natural measure of geometric **similarity**

What is Optimal Transport?

A **geometric** way
to compare
probability measures.

Nobel prize



Monge



Kantorovich



Dantzig



Wasserstein



Brenier



Otto



McCann



Villani

Fields medal
(and French politician)



Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

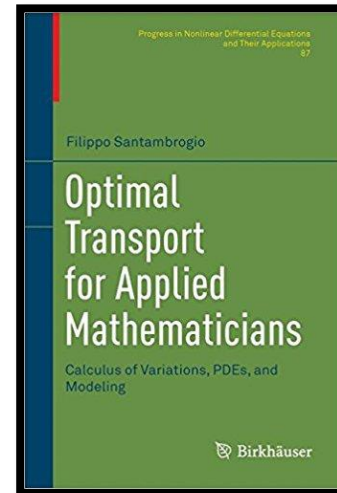
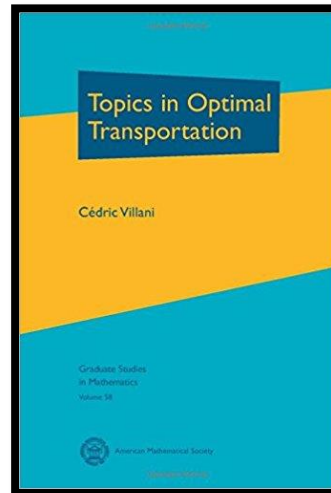
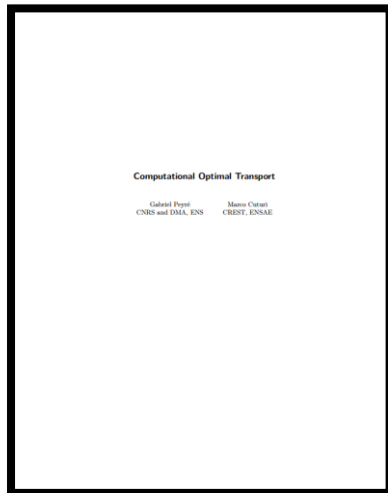
2. Applications

3. Discrete/discretized transport

- Entropic regularization
- Eulerian transport
- Semidiscrete transport

4. Extensions & frontiers

Useful References



Snapshots of modern mathematics
from Oberwolfach

N° 8/2017

Computational Optimal Transport

Justin Solomon

Optimal transport is the mathematical discipline of matching supply to demand while minimizing shipping costs. This matching problem becomes extremely challenging as the quantity of supply and demand points increases; modern applications must cope with thousands or millions of these at a time. Here, we introduce the computational optimal transport prob-

Optimal Transport on Discrete Domains

Justin Solomon

ABSTRACT. Inspired by the matching of supply to demand in logistical problems, the optimal transport (or Monge–Kantorovich) problem involves the matching of probability distributions defined over a geometric domain such as a surface or manifold. In its most obvious discretization, optimal transport becomes a large-scale linear program, which typically is infeasible to solve efficiently on triangle meshes, graphs, point clouds, and other domains encountered in graphics and machine learning. Recent breakthroughs in numerical optimal transport, however, enable scalability to orders-of-magnitude larger problems, solvable in a fraction of a second. Here, we discuss advances in numerical optimal transport that leverage understanding of both discrete and smooth aspects of the problem. State-of-the-art techniques in discrete optimal transport combine insight from partial differential equations (PDE) with convex analysis to reformulate, discretize, and optimize transportation problems. The end result is a set of theoretically-justified models suitable for domains with thousands or millions of vertices. Since numerical optimal transport is a relatively new discipline, special emphasis is placed on identifying and explaining open problems in need of mathematical insight and additional research.

1. Introduction

Many tools from discrete differential geometry (DDG) were inspired by practical considerations in areas like computer graphics and vision. Disciplines like these require fine-grained understanding of geometric structure and the relationships between different shapes—problems for which the toolbox from smooth geometry can provide substantial insight. Indeed, a triumph of discrete differential geometry is its

Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

2. Applications

3. Discrete/discretized transport

- Entropic regularization
- Eulerian transport
- Semidiscrete transport

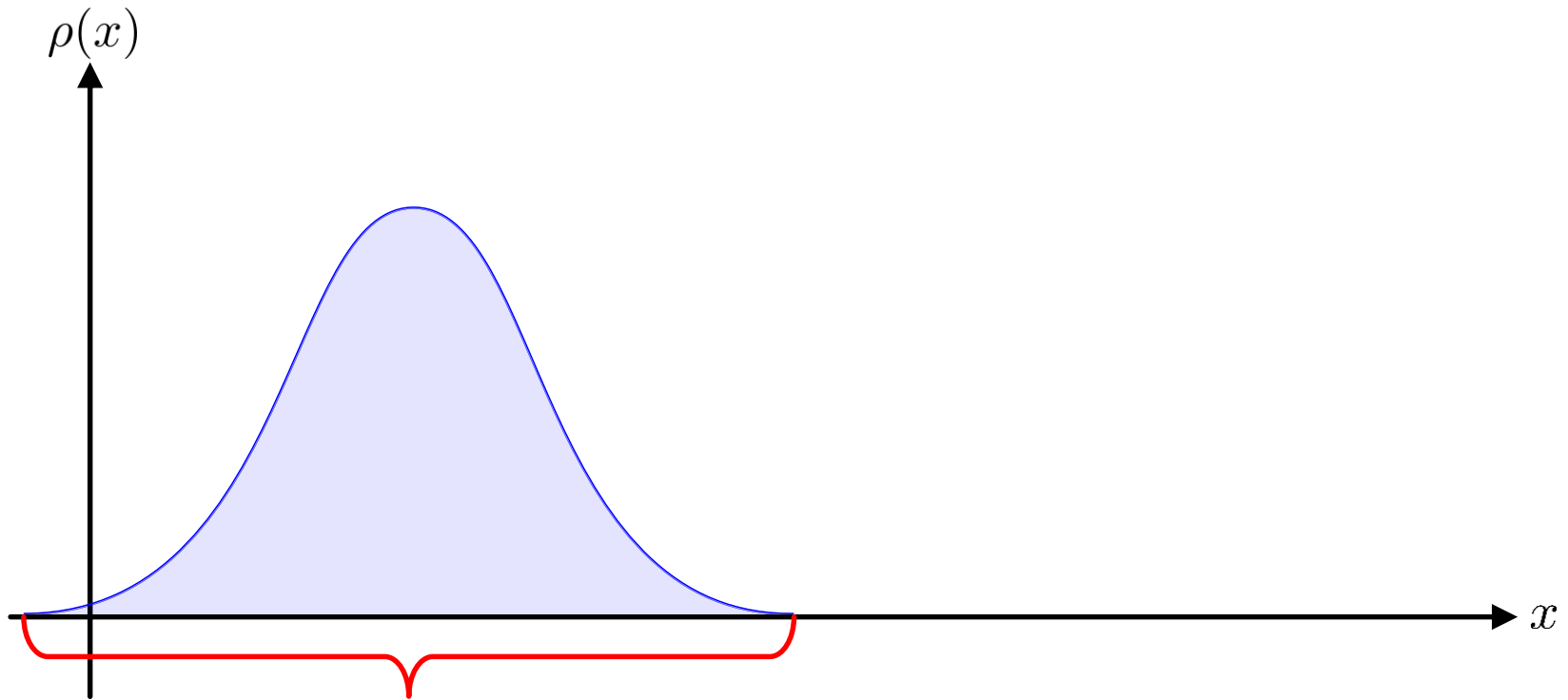
4. Extensions & frontiers



Transport Philosophy

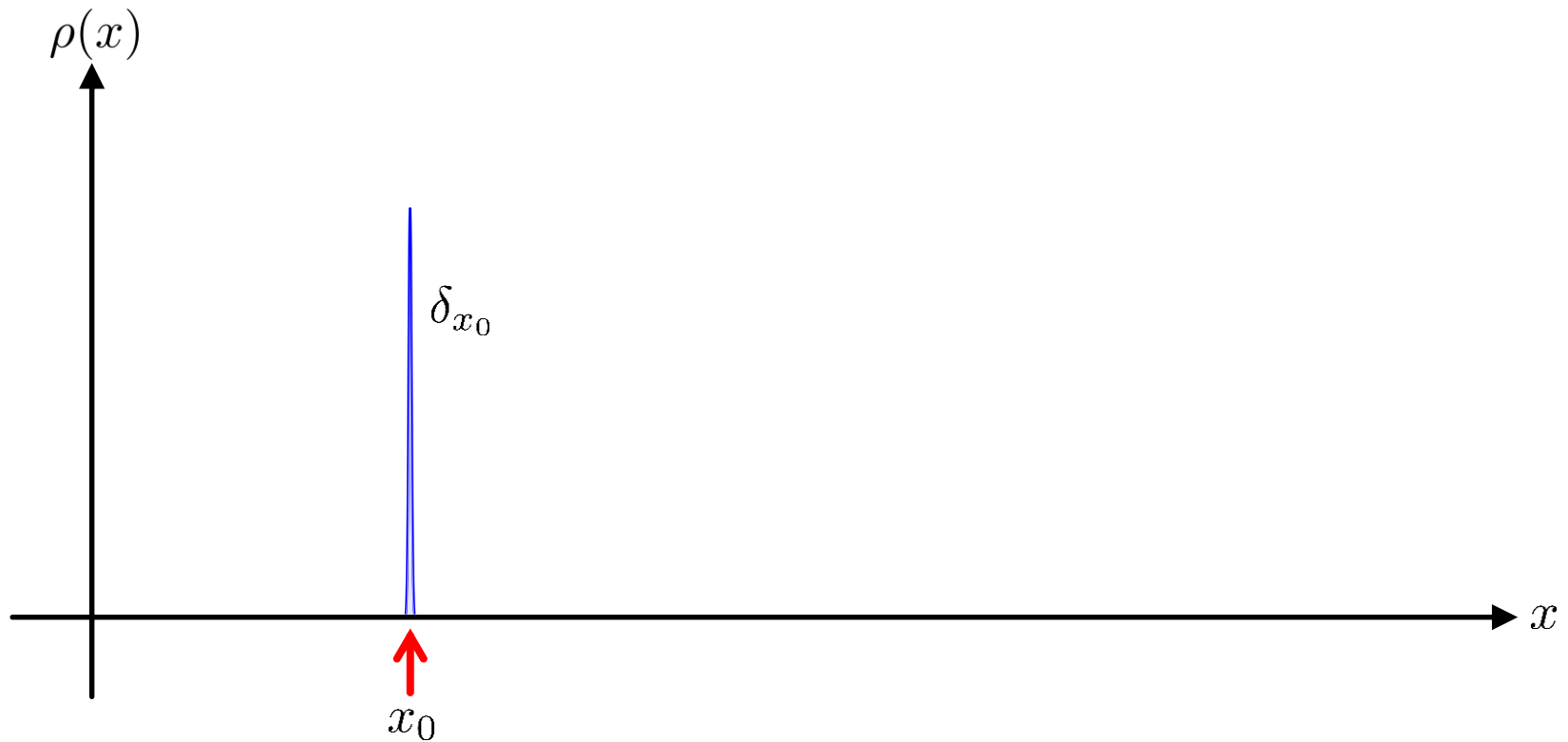
Understand geometry from a
“softened” probabilistic
standpoint.

Probability as Geometry



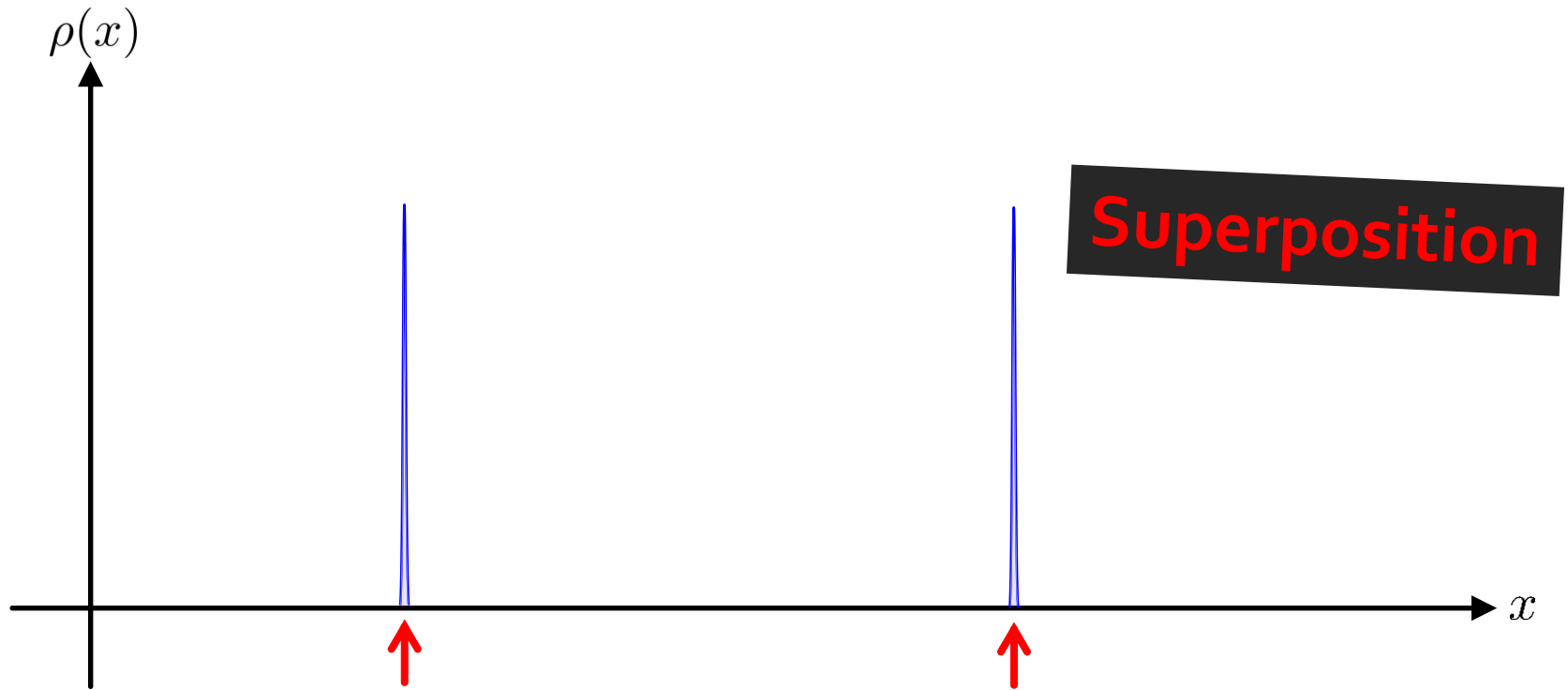
“Somewhere over here.”

Probability as Geometry



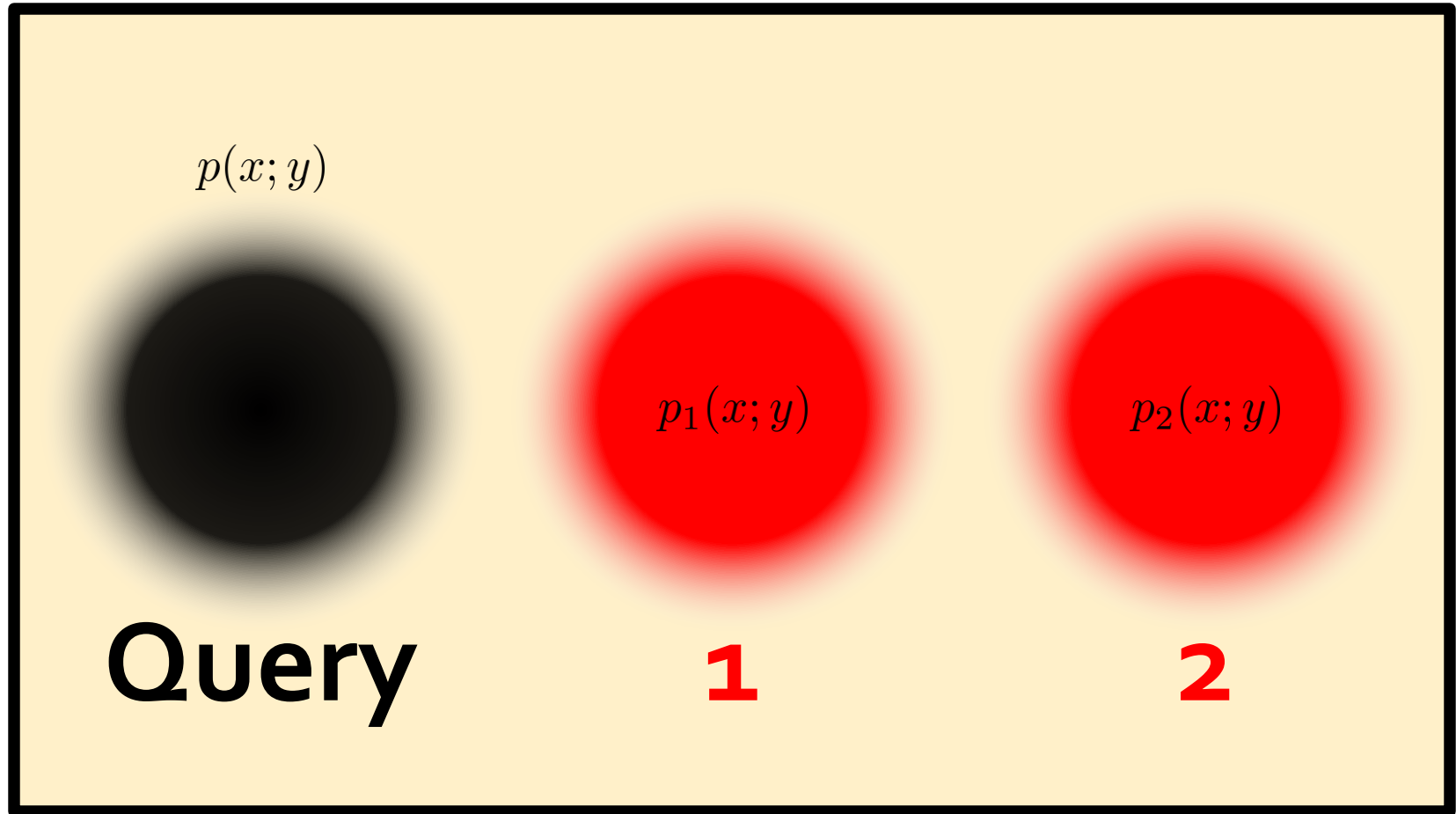
“Exactly here.”

Probability as Geometry



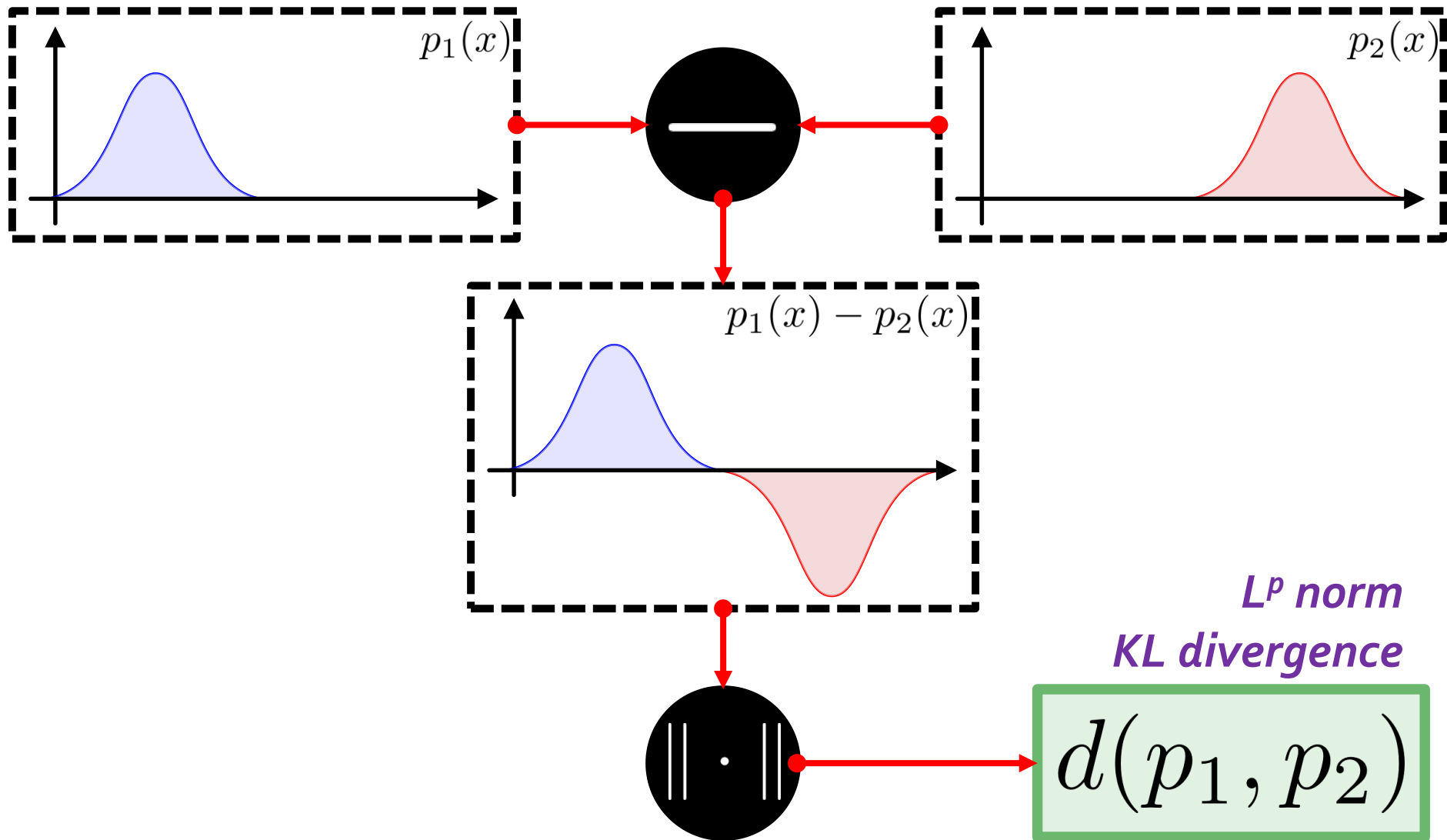
“One of these two places.”

Fuzzy Geometry

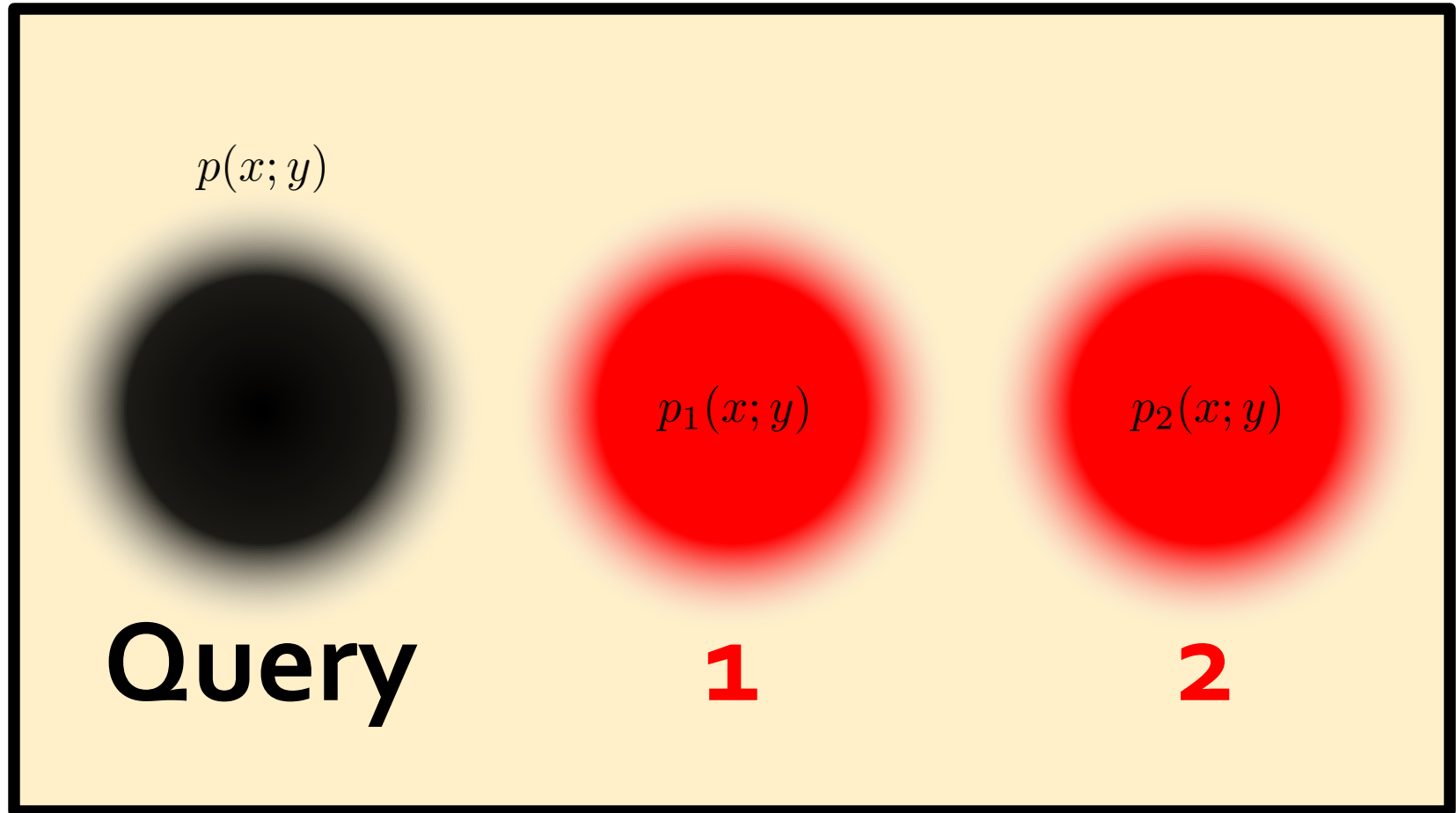


Which is closer, 1 or 2?

Typical Measurement

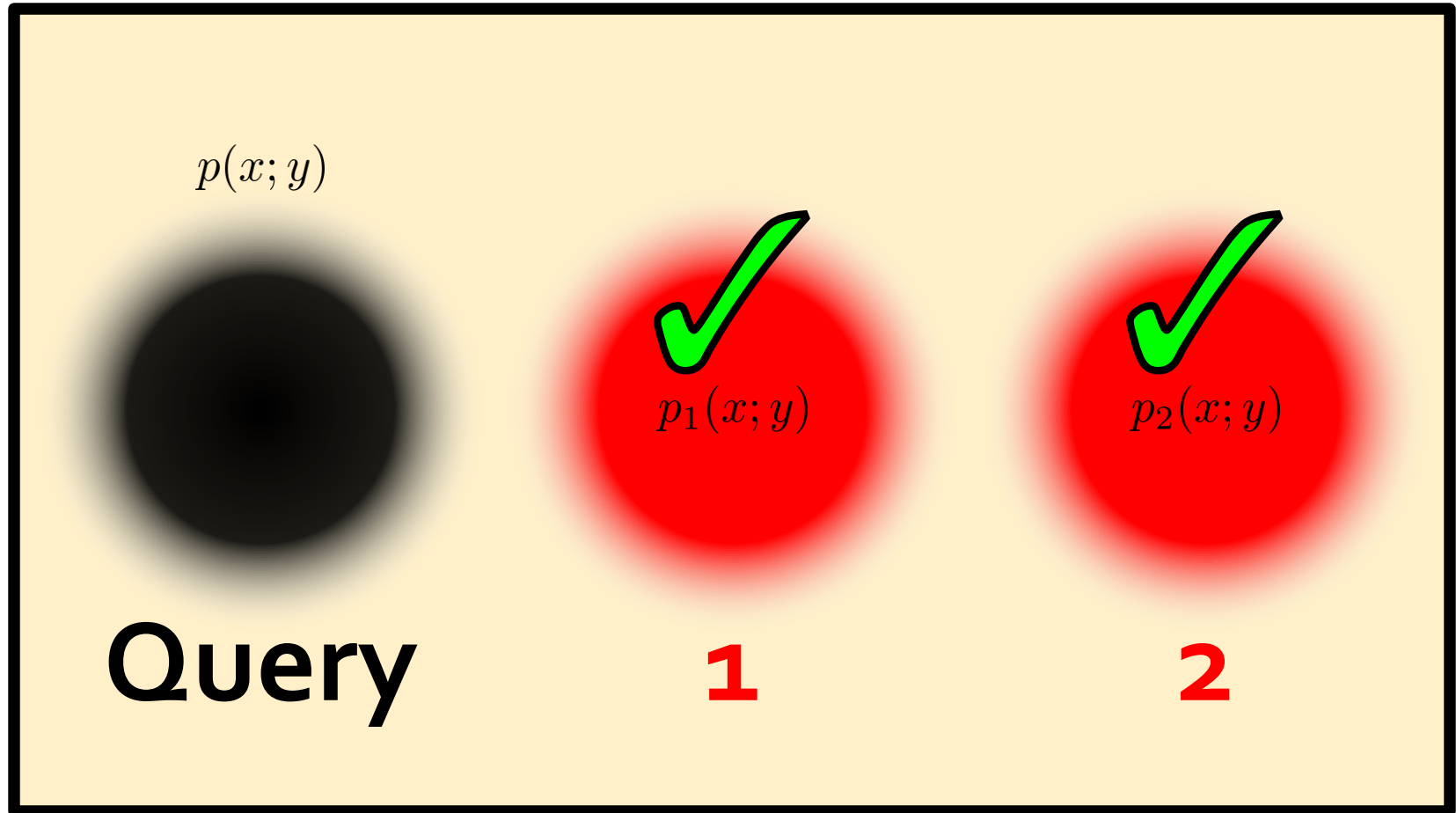


Returning to the Question



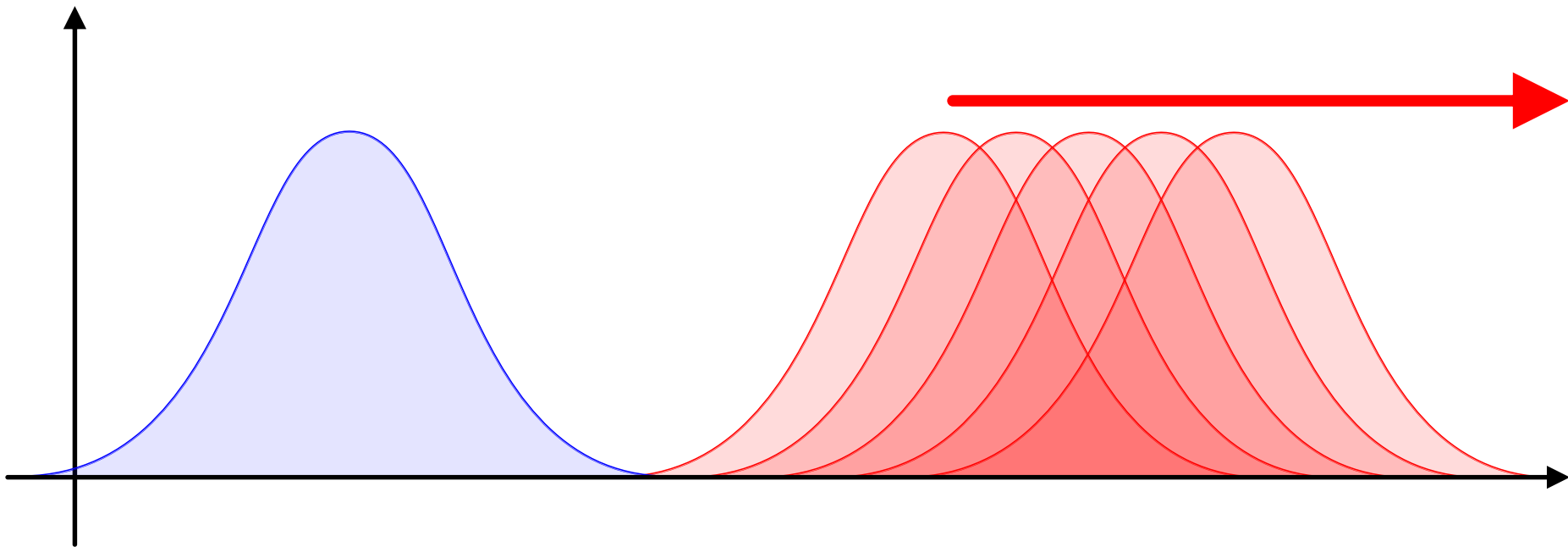
Which is closer, 1 or 2?

Returning to the Question



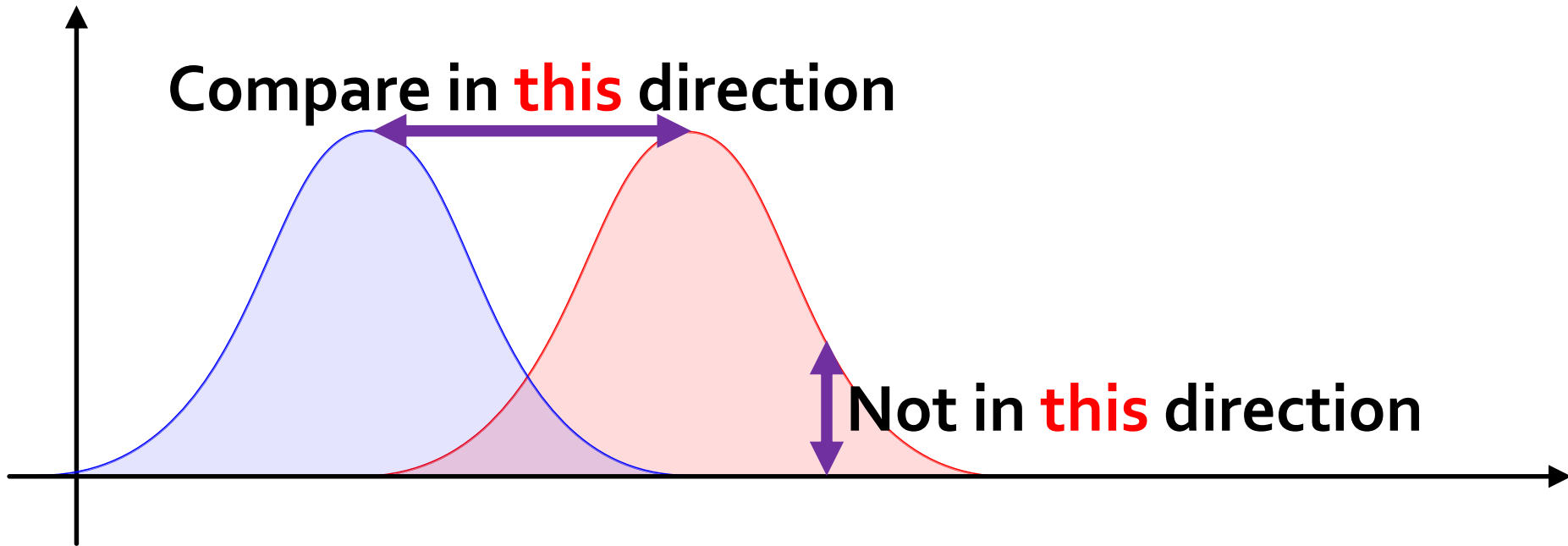
Neither! Equidistant.

What's Wrong?



**Measured overlap,
not displacement.**

Alternative Idea



Observation

Even the laziest shoveler
must do some work.

Property of the distributions themselves!



My house!

The Setup: Transport in 1D

$\pi(x, y) :=$ Amount moved from x to y

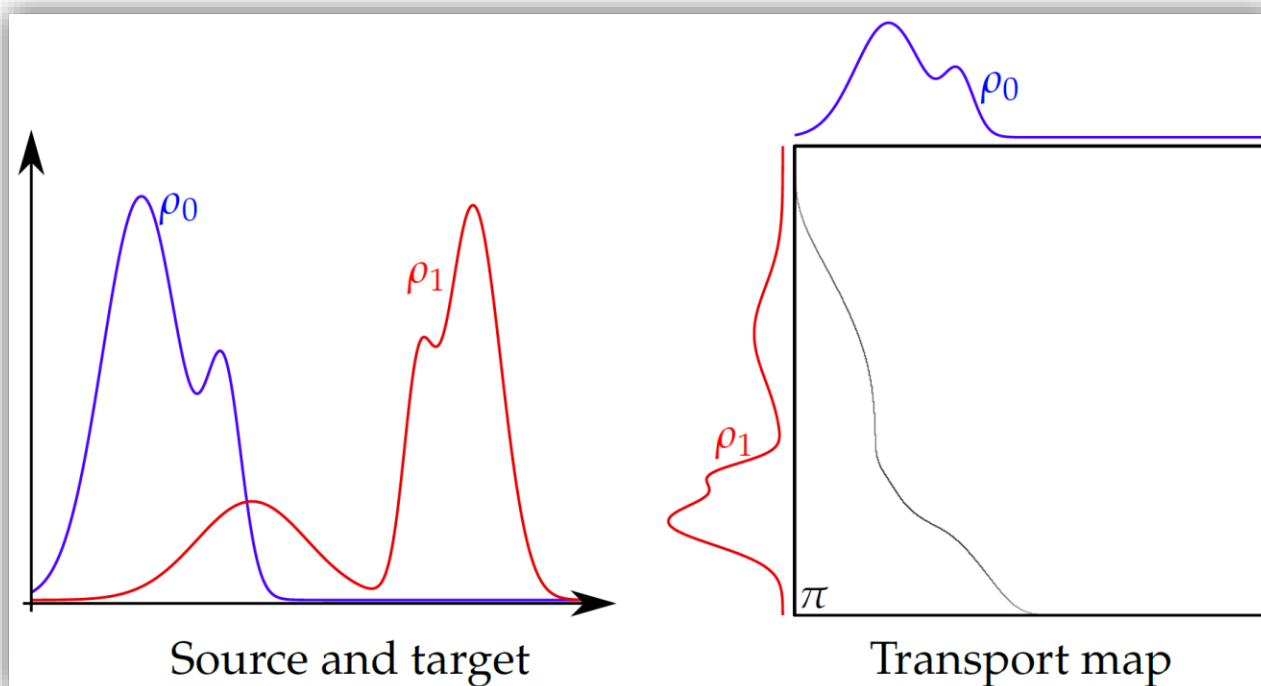
$$\pi(x, y) \geq 0 \quad \forall x, y \in \mathbb{R} \quad \text{Mass is positive}$$

$$\int_{\mathbb{R}} \pi(x, y) dy = \rho_0(x) \quad \forall x \in \mathbb{R} \quad \text{Must scoop everything up}$$

$$\int_{\mathbb{R}} \pi(x, y) dx = \rho_1(y) \quad \forall y \in \mathbb{R} \quad \text{Must cover the target}$$

1-Wasserstein in 1D

$$\mathcal{W}_1(\rho_0, \rho_1) := \begin{cases} \min_{\pi} & \iint_{\mathbb{R} \times \mathbb{R}} \pi(x, y) |x - y| dx dy & \text{Minimize total work} \\ \text{s.t.} & \pi \geq 0 \ \forall x, y \in \mathbb{R} & \text{Nonnegative mass} \\ & \int_{\mathbb{R}} \pi(x, y) dy = \rho_0(x) \ \forall x \in \mathbb{R} & \text{Starts from } \rho_0 \\ & \int_{\mathbb{R}} \pi(x, y) dx = \rho_1(y) \ \forall y \in \mathbb{R} & \text{Ends at } \rho_1 \end{cases}$$



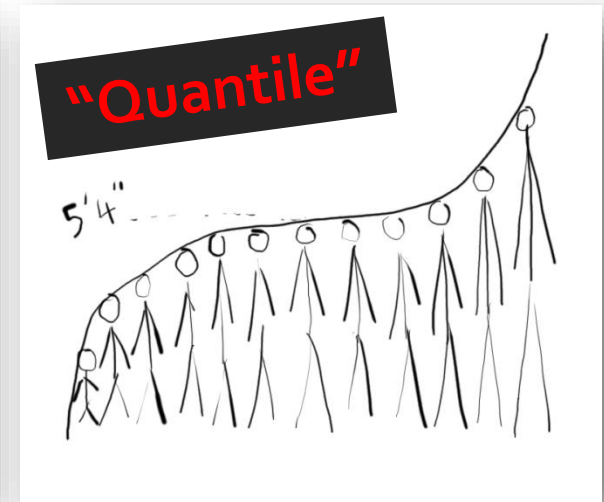
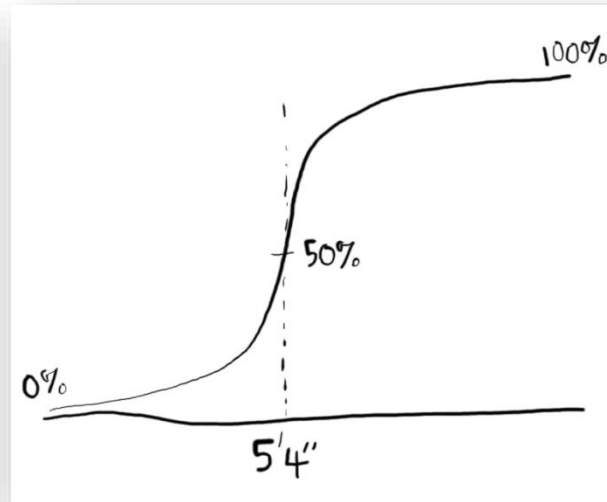
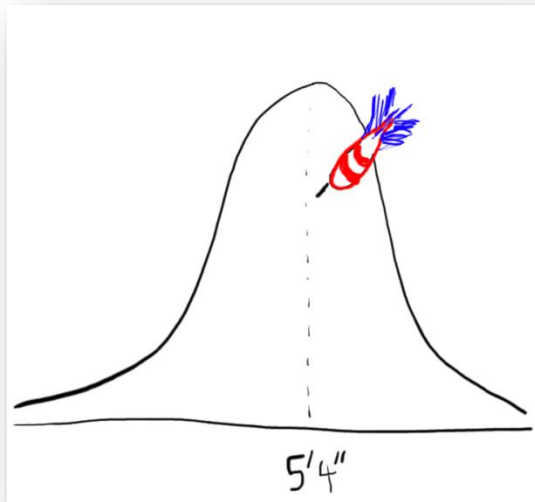


When is transport **computable?**

Needed: Finite number of unknowns.

In One Dimension: Closed-Form

<http://realgl.blogspot.com/2013/01/pdf-cdf-inv-cdf.html>



PDF **[CDF]** **CDF⁻¹**

$$\mathcal{W}_1(\mu, \nu) = \int_{-\infty}^{\infty} |\text{CDF}(\mu) - \text{CDF}(\nu)| d\ell$$

$$\mathcal{W}_2^2(\mu, \nu) = \int_{-\infty}^{\infty} (\text{CDF}^{-1}(\mu) - \text{CDF}^{-1}(\nu))^2 d\ell$$

In One Dimension: Closed-Form

<http://realgl.blogspot.com/2013/01/pdf-cdf-inv-cdf.html>



**Doesn't extend
past 1D!**

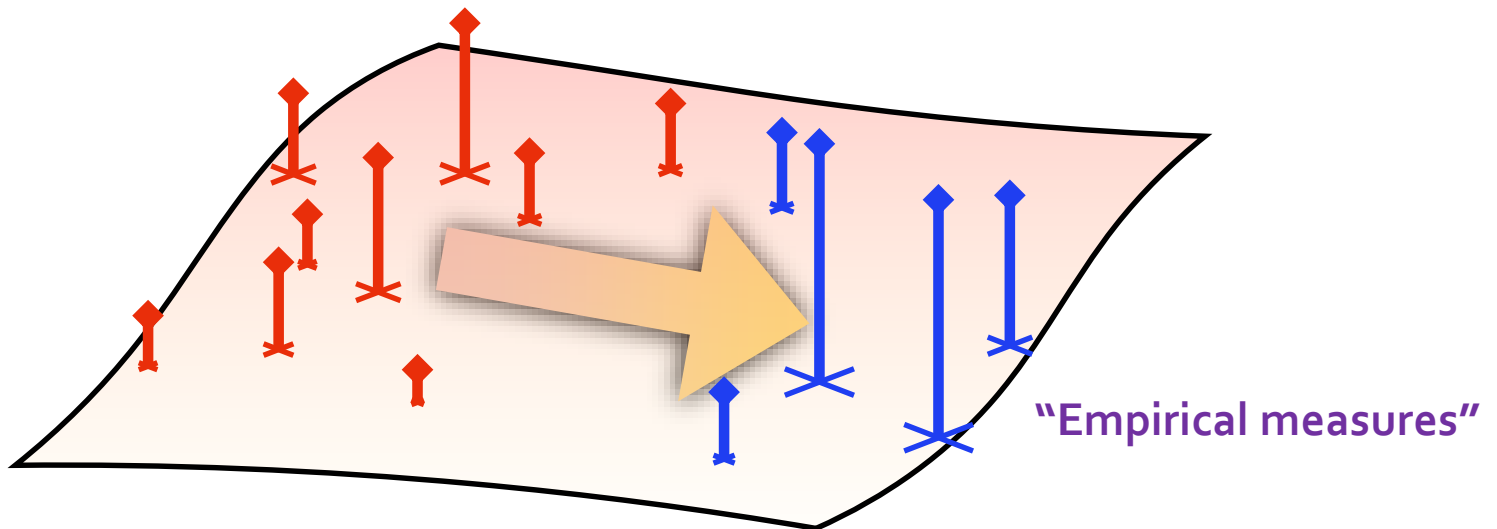
$$W_2^2(\mu, \nu) = \int_{-\infty}^{\infty} |CDF(\mu) - CDF(\nu)| d\ell$$
$$W_2^2(\mu, \nu) = \int_{-\infty}^{\infty} (CDF^{-1}(\mu) - CDF^{-1}(\nu))^2 d\ell$$

Fully-Discrete Transport

$$[\mathcal{W}_p(\mu_0, \mu_1)]^p = \begin{cases} \min_{T \in \mathbb{R}^{k_0 \times k_1}} & \sum_{ij} T_{ij} |x_{0i} - x_{1j}|^p \\ \text{s.t.} & T \geq 0 \\ & \sum_j T_{ij} = a_{0i} \\ & \sum_i T_{ij} = a_{1j} \end{cases}$$

Linear program: Finite number of variables

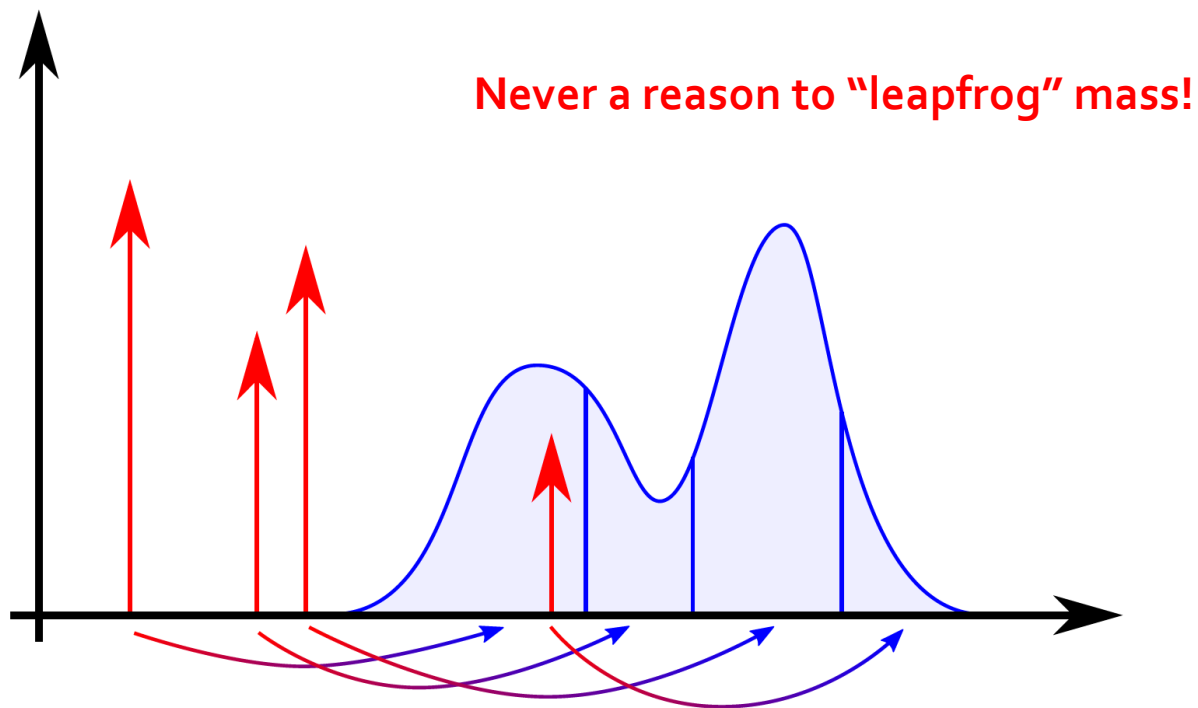
Algorithms: Simplex, interior point, auction, ...



Semidiscrete Transport

$$\mu_0 := \sum_{i=1}^{k_0} a_{0i} \delta_{x_{0i}}$$

$$\mu_1(S) := \int_S \rho_1(x) dx$$



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Probability Measure

$$\mu(X) = 1$$

$$\mu(S \subseteq X) \in [0, 1]$$

X is the
domain

$$\mu\left(\bigcup_{i \in I} E_i\right) = \sum_{i \in I} \mu(E_i)$$

“Prob(X)”

when E_i disjoint,
 I countable

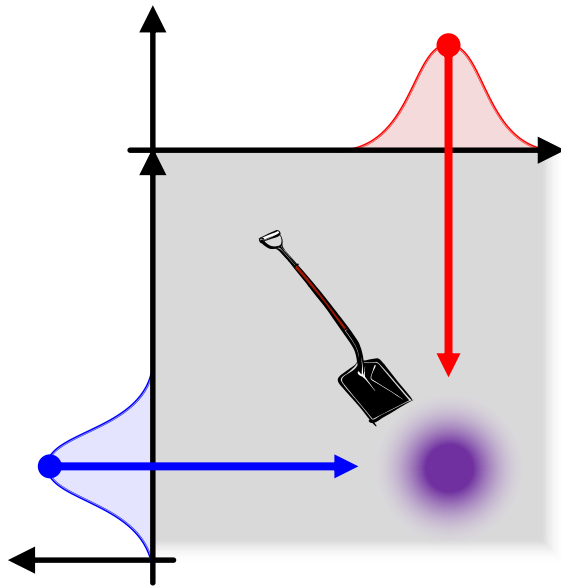
Function from sets to probability

Measure Coupling

$$\mu \in \text{Prob}(X), \nu \in \text{Prob}(Y)$$



$$\Pi(\mu, \nu) := \left\{ \pi \in \text{Prob}(X \times Y) : \begin{pmatrix} \pi(A \times Y) = \mu(A) \\ \pi(X \times B) = \nu(B) \end{pmatrix} \right\}$$



**Analog of
transportation
matrix**

Kantorovich Problem

$$\text{OT}(\mu, \nu; c) := \min_{\pi \in \Pi(\mu, \nu)} \iint_{X \times Y} c(x, y) d\pi(x, y)$$

General transport problem!

Example: Discrete Transport

$$X = \{1, 2, \dots, k_1\}, Y = \{1, 2, \dots, k_2\}$$

$$\text{OT}(v, w; C) = \begin{cases} \min_{T \in \mathbb{R}^{k_1 \times k_2}} & \sum_{ij} T_{ij} c_{ij} \\ \text{s.t.} & T \geq 0 \\ & \sum_j T_{ij} = v_i \quad \forall i \in \{1, \dots, k_1\} \\ & \sum_i T_{ij} = w_j \quad \forall j \in \{1, \dots, k_2\}. \end{cases} \quad \text{"Earth Mover's Distance"}$$

Metric when $d(x, y)$ satisfies the triangle inequality.

"The Earth Mover's Distance as a Metric for Image Retrieval"

Rubner, Tomasi, and Guibas; IJCV 40.2 (2000): 99—121.

Revised in:

"Ground Metric Learning"

Cuturi and Avis; JMLR 15 (2014)

p -Wasserstein Distance

$$\mathcal{W}_p(\mu, \nu) \equiv \min_{\pi \in \Pi(\mu, \nu)} \left(\iint_{X \times X} \underset{\substack{\uparrow \\ \text{Shortest path} \\ \text{distance}}}{d(x, y)^p} d\pi(x, y) \right)^{1/p}$$

Expectation



Kantorovich Duality

$$\begin{aligned} \text{OT}(\mu, \nu; c) &:= \begin{cases} \min_{\pi} \iint_{X \times Y} c(x, y) d\pi(x, y) \\ \text{s.t. } \pi \in \Pi(\mu, \nu) \end{cases} && \text{Primal} \\ &= \begin{cases} \max_{\phi, \psi} \int_X \phi(x) d\mu(x) + \int_Y \psi(y) d\nu(y) \\ \text{s.t. } \phi(x) + \psi(y) \leq c(x, y) \text{ for a.e. } x \in X, y \in Y \end{cases} && \text{Dual} \end{aligned}$$

Flow-Based W_2

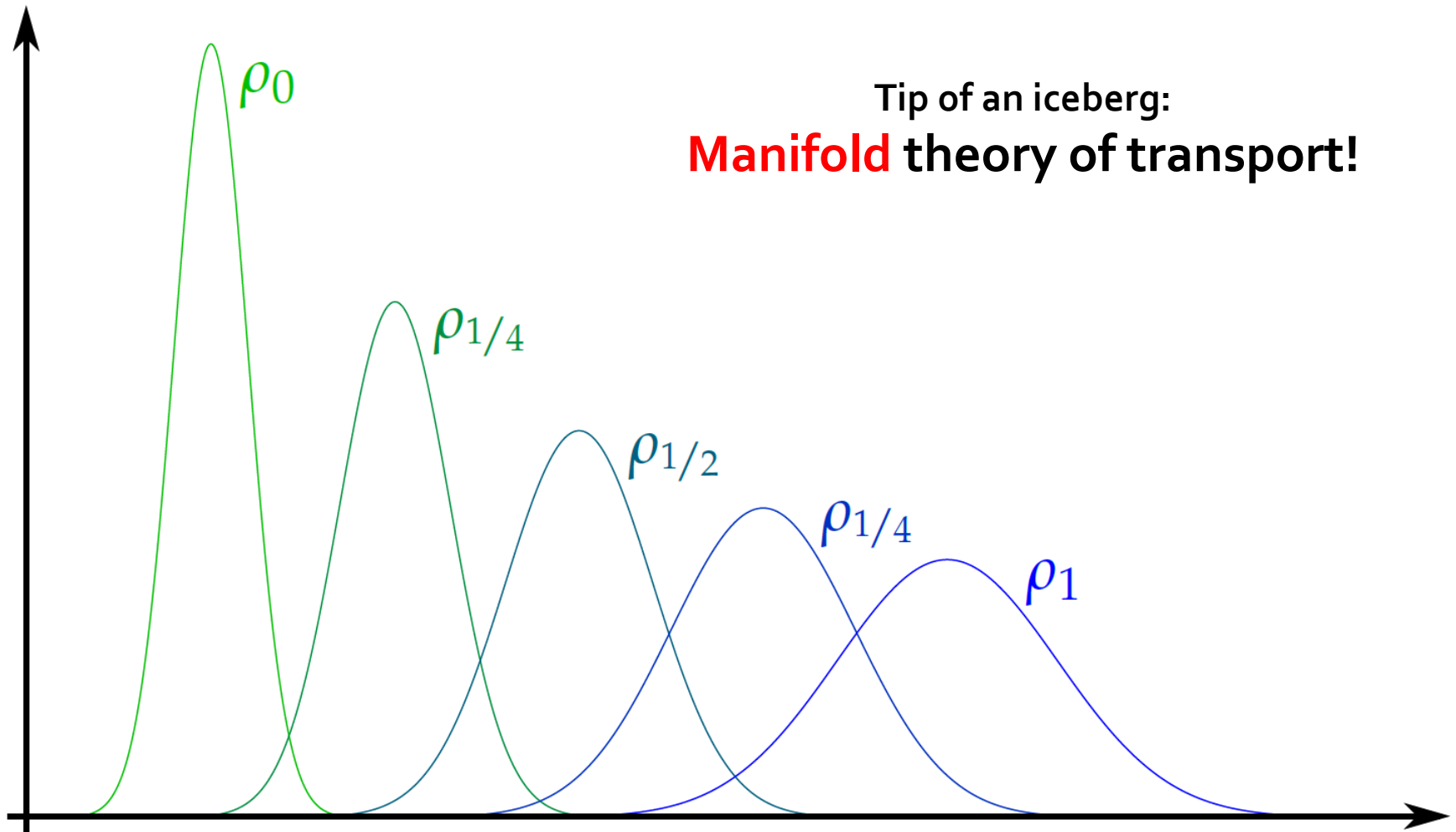
$$\mathcal{W}_2^2(\rho_0, \rho_1) = \begin{cases} \inf_{\rho, v} \iint_{M \times [0,1]} \frac{1}{2} \rho(x, t) \|v(x, t)\|^2 dx dt \\ \text{s.t. } \nabla \cdot (\rho(x, t) v(x, t)) = \frac{\partial \rho(x, t)}{\partial t} \\ v(x, t) \cdot \hat{n}(x) = 0 \quad \forall x \in \partial M \\ \rho(x, 0) = \rho_0(x) \\ \rho(x, 1) = \rho_1(x) \end{cases}$$

Benamou & Brenier

“A computational fluid mechanics solution of the Monge-Kantorovich mass transfer problem”

Numer. Math. 84 (2000), pp. 375-393

Displacement Interpolation



Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

2. Applications

3. Discrete/discretized transport

- Entropic regularization
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- Semidiscrete transport

4. Extensions & frontiers



Wassersteinization

[wos-ur-stahyn-ahy-sey-shuh-n]

noun.

Introduction of optimal transport
into a computational problem.

cf. least-squarification, L_1 ification, deep-netification, kernelization

Key Ingredients

We have tools to

- **Solve** optimal transport problems numerically
- **Differentiate** transport distances in terms of their input distributions

Bonus:

Transport cost from μ to ν is a **convex** function of μ and ν .

Redistricting?

Balanced power diagrams for redistricting

Vincent Cohen-Addad*

Philip N. Klein[†]

Neal E. Young[‡]

January 6, 2018

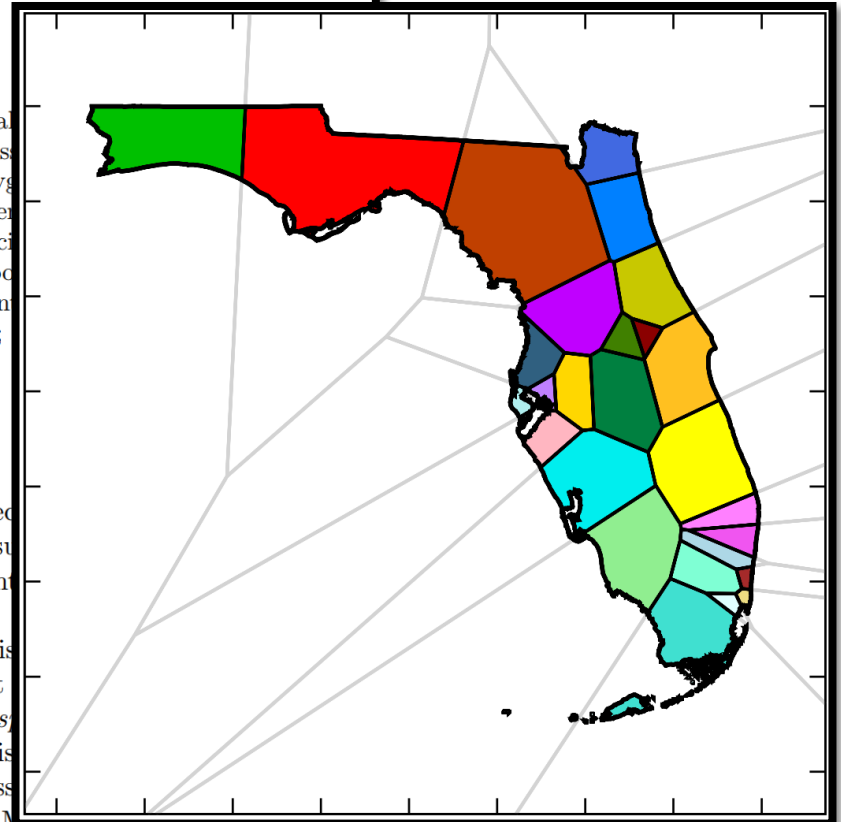
Abstract

We explore a method for *redistricting*, decomposing a geographical area into *districts*, so that the populations of the districts are as close as possible to equal, compact and contiguous. Each district is the intersection of a polygon and a Voronoi cell. The polygons are convex and the average number of sides per district is small. The polygons tend to be quite compact. With each polygon is associated a centroid, the centroid of the locations of the residents associated with the polygon. This can be viewed as a heuristic for finding centers and a balanced assignment of residents to centers, as to minimize the sum of squared distances of residents to centers; this method is said to have low dispersion.

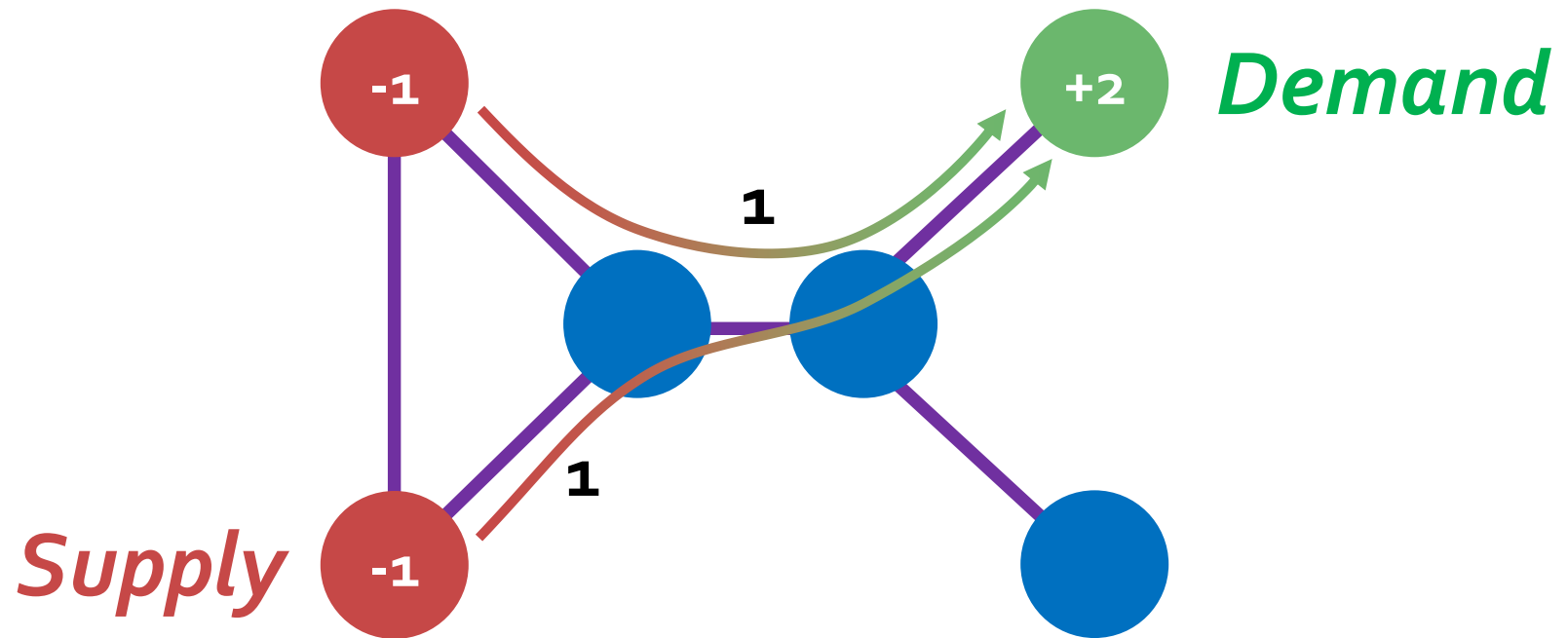
1 Introduction

Redistricting. *Redistricting*, in the context of elections refers to dividing a geographical area into subareas such that all subareas have the same population. The subareas are supposed to be *contiguous* to the extent that the subarea can reasonably be interpreted to mean *connected*.

In most states, districts are also supposed to be *compact*. This is a difficult concept to define. Some measures of compactness are based on boundaries; a district is compact if its boundary is simpler rather than contorted. Some measures are based on *dispersion*, the distance the district spreads from a central core” [17]. Idaho directs its redistricting commission to “draw districts that are oddly shaped.” Other states loosely address the issue. “Arizona and Colorado focus on contorted boundaries; California, Massachusetts, and Minnesota focus on dispersion; and Iowa embraces both” [17].

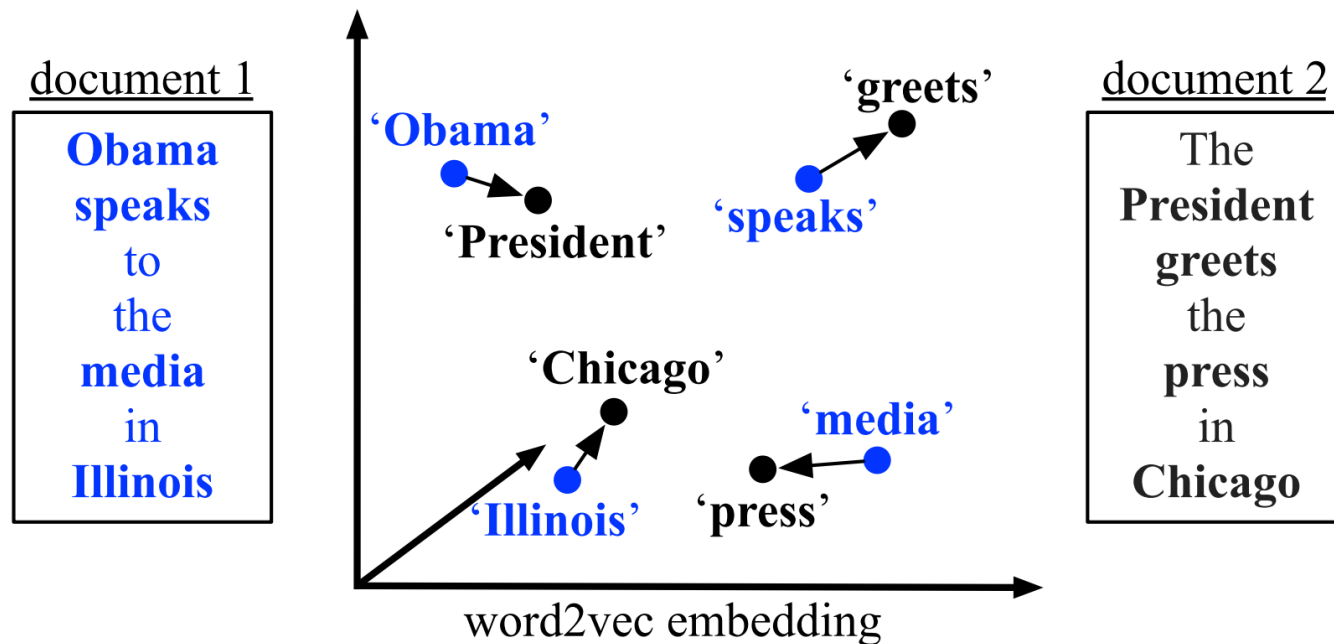


Operations and Logistics



Minimum-cost flow

Histograms and Descriptors



Use deep network embedding

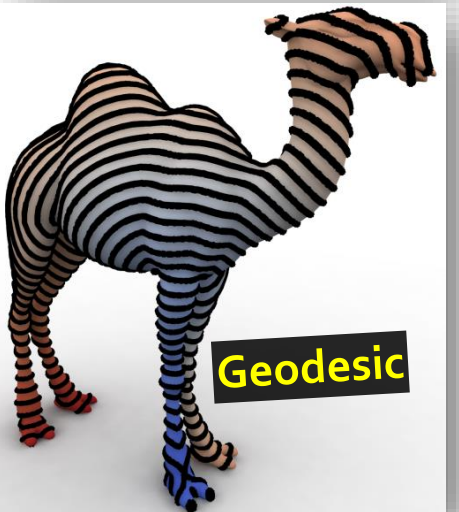
[Kusner et al. 2015]

Word Mover's Distance (WMD)

Distance Approximation



Biharmonic



Geodesic



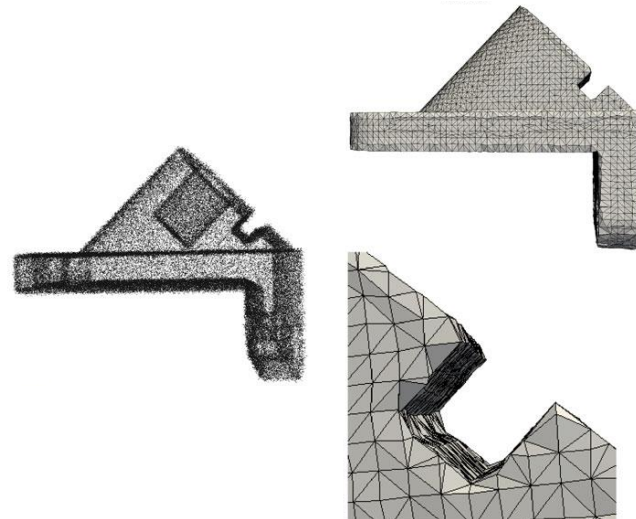
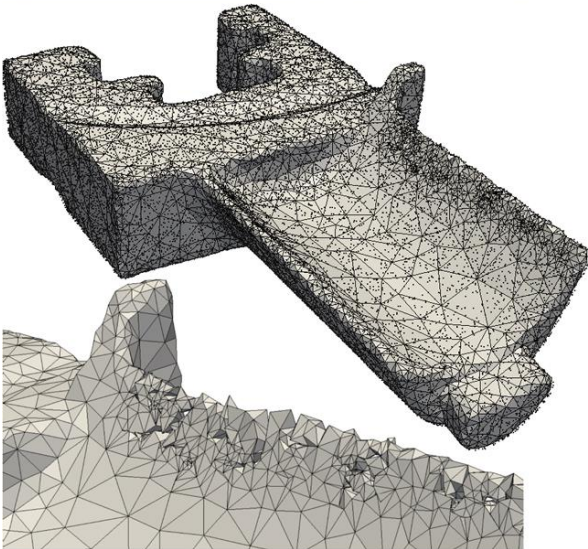
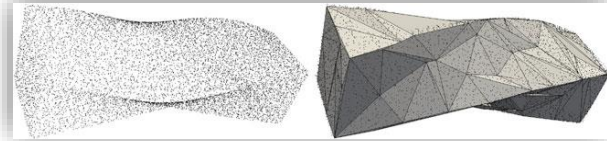
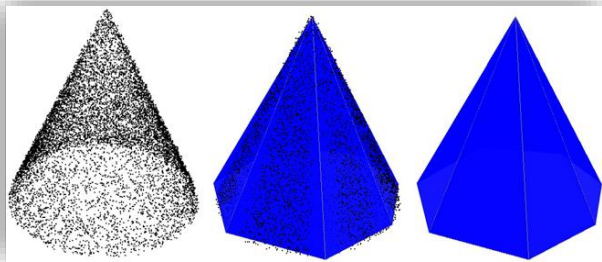
0 eigenfunctions



100 eigenfunctions

Proposition: Satisfies triangle inequality.

Registration and Reconstruction



[Digne et al. 2014]

Distance from point cloud to mesh

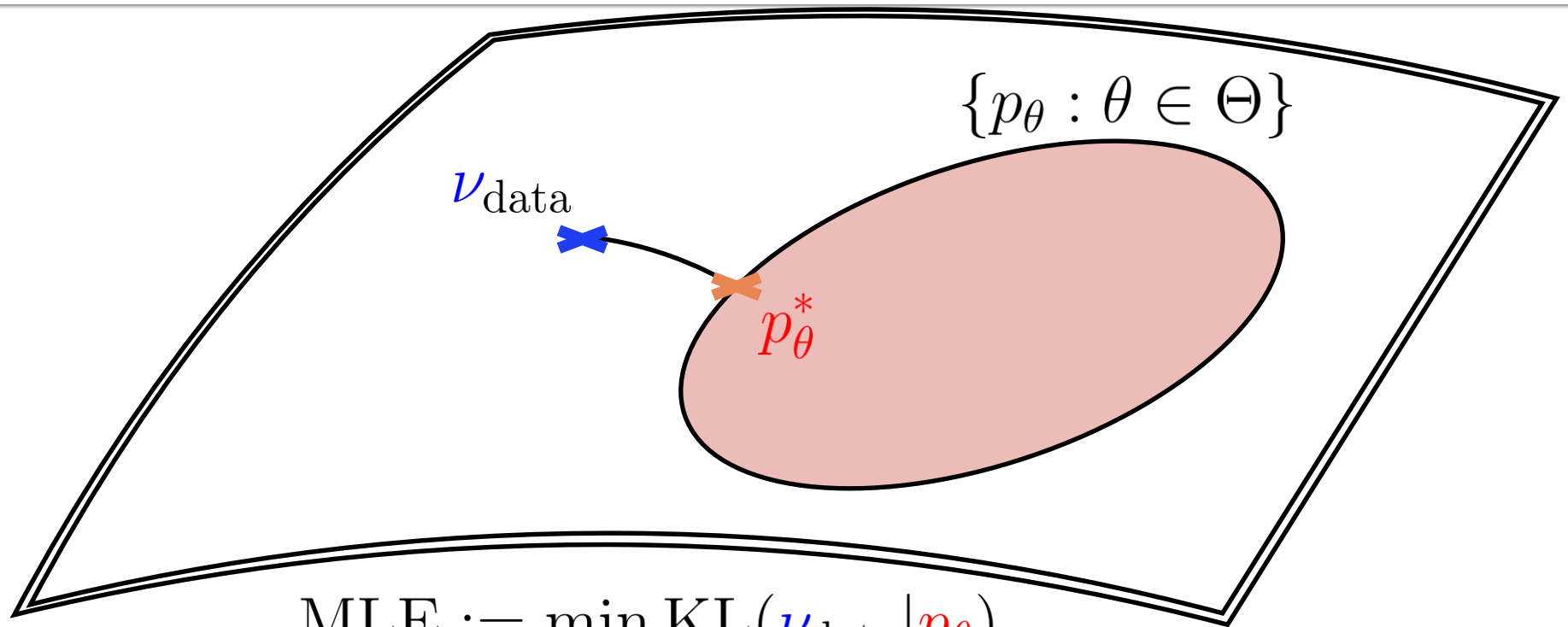
Blue Noise and Stippling



$$\min_{x_1, \dots, x_n} \mathcal{W}_2^2 \left(\mu, \frac{1}{n} \sum_i \delta_{x_i} \right)$$

Image courtesy F. de Goes; photo by F. Durand

Statistical Estimation



$$\text{MLE} := \min_{\theta \in \Theta} \text{KL}(\nu_{\text{data}} | p_\theta)$$

$$\longrightarrow \text{MKE} := \min_{\theta \in \Theta} \mathcal{W}_2(\nu_{\text{data}}, p_\theta)$$

[Bassetti 2006]

Minimum Kantovich Estimator

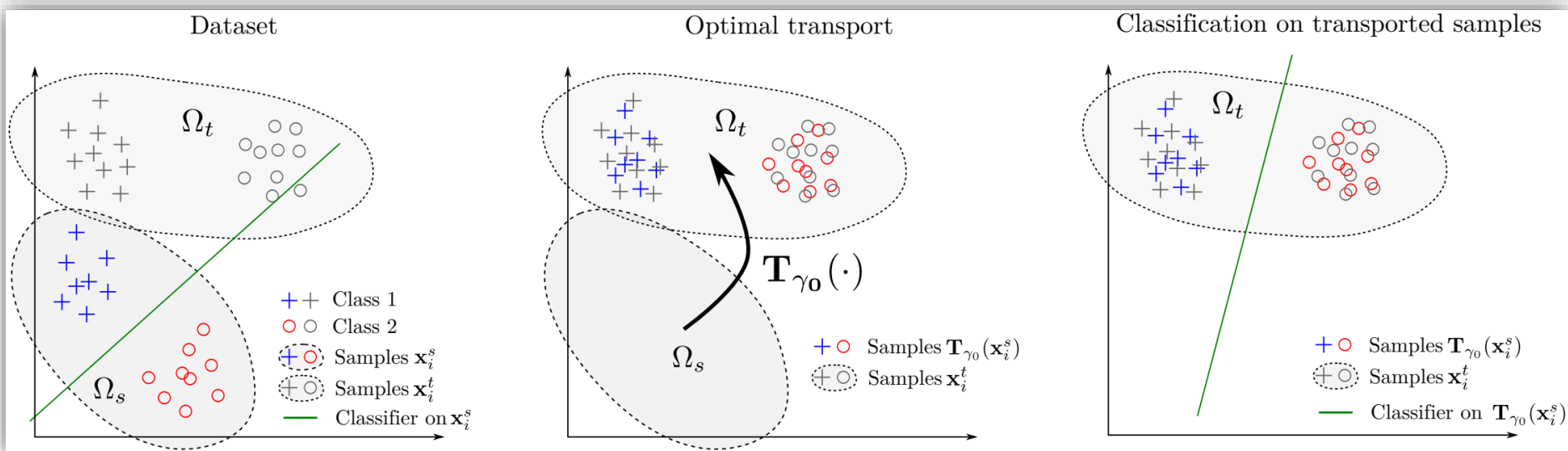
Distributionally Robust Optimization

$$\inf_{x \in \mathbb{X}} \sup_{Q \in \hat{\mathcal{P}}_N} \mathbb{E}_{\xi \sim Q} [h(x, \xi)]$$

**Wasserstein ball around
empirical distribution**

Loss function

Domain Adaptation



1. Estimate transport map
2. Transport labeled samples to new domain
3. Train classifier on transported labeled samples

Engineering Design



Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

2. Applications

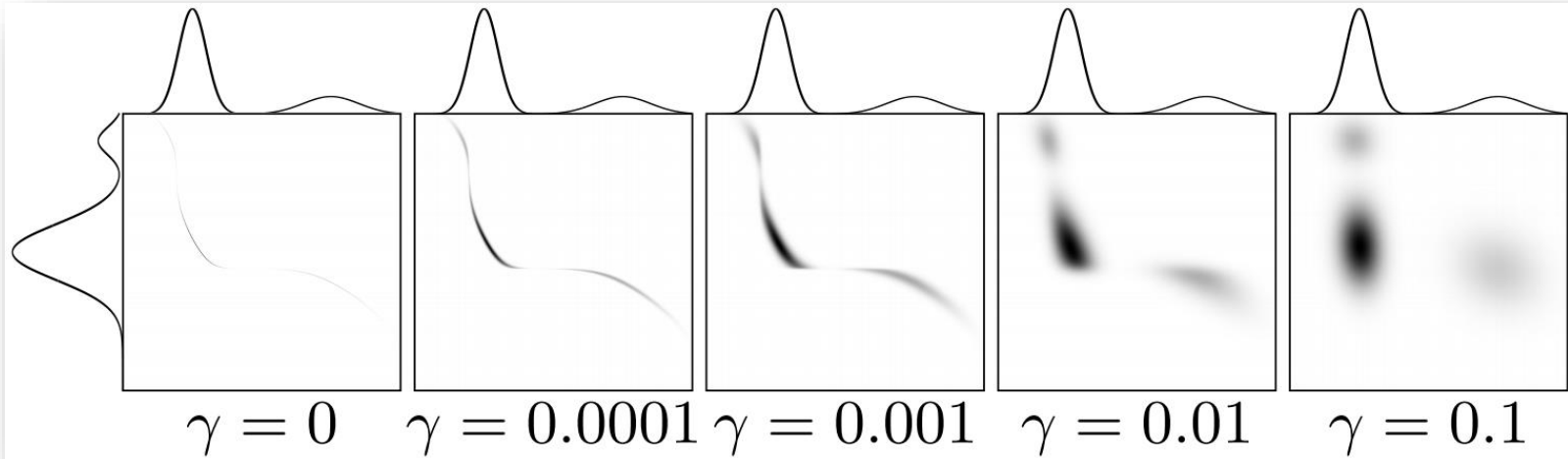
3. Discrete/discretized transport

- Entropic regularization
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Entropic Regularization



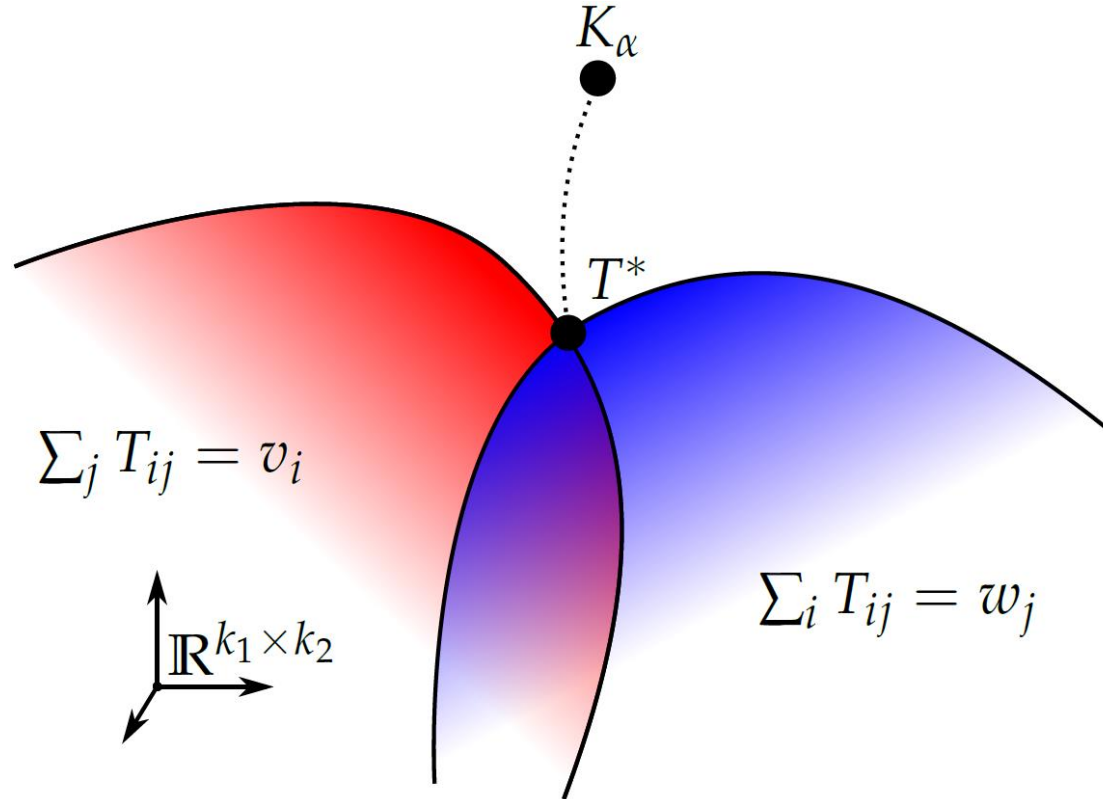
$$\begin{aligned} \min_T \quad & \sum_{ij} T_{ij} c_{ij} - \alpha H(T) \\ \text{s.t.} \quad & \sum_j T_{ij} = v_i \\ & \sum_i T_{ij} = w_j \end{aligned}$$

OK to drop
nonnegative
constraint!

$$H(T) := - \sum_{ij} T_{ij} \log T_{ij}$$

Interpretation as Projection

$$\sum_{ij} T_{ij} c_{ij} - \alpha H(T) = \text{KL}(T|K_\alpha) \text{ where } K_\alpha := \exp(-C_\alpha)$$

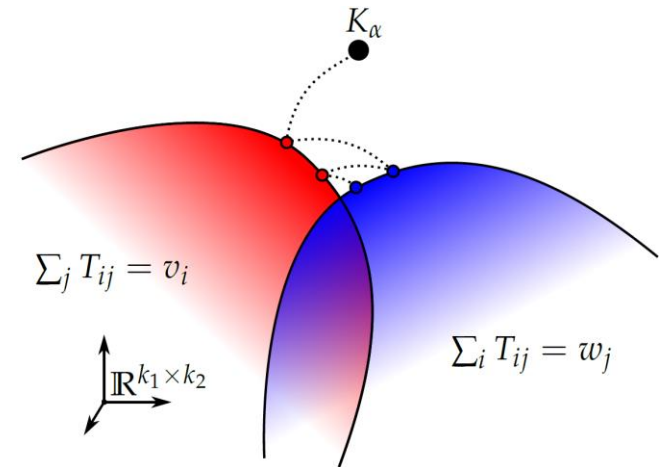


Sinkhorn Algorithm

$$T = \text{diag}(p) K_\alpha \text{diag}(q),$$

where $K_\alpha := \exp(-C/\alpha)$

$$p \leftarrow v \oslash (K_\alpha q)$$
$$q \leftarrow w \oslash (K_\alpha^\top p)$$

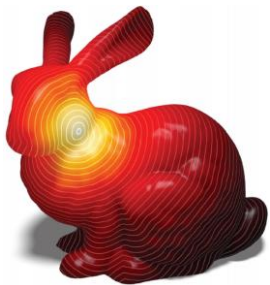


Sinkhorn & Knopp. "Concerning nonnegative matrices and doubly stochastic matrices".
Pacific J. Math. 21, 343–348 (1967).

Alternating projection

Ingredients for Sinkhorn

1. Supply vector p
2. Demand vector q
3. **Multiplication** by K



$$K_{ij} = e^{-c_{ij}/\alpha}$$

Plan For Today

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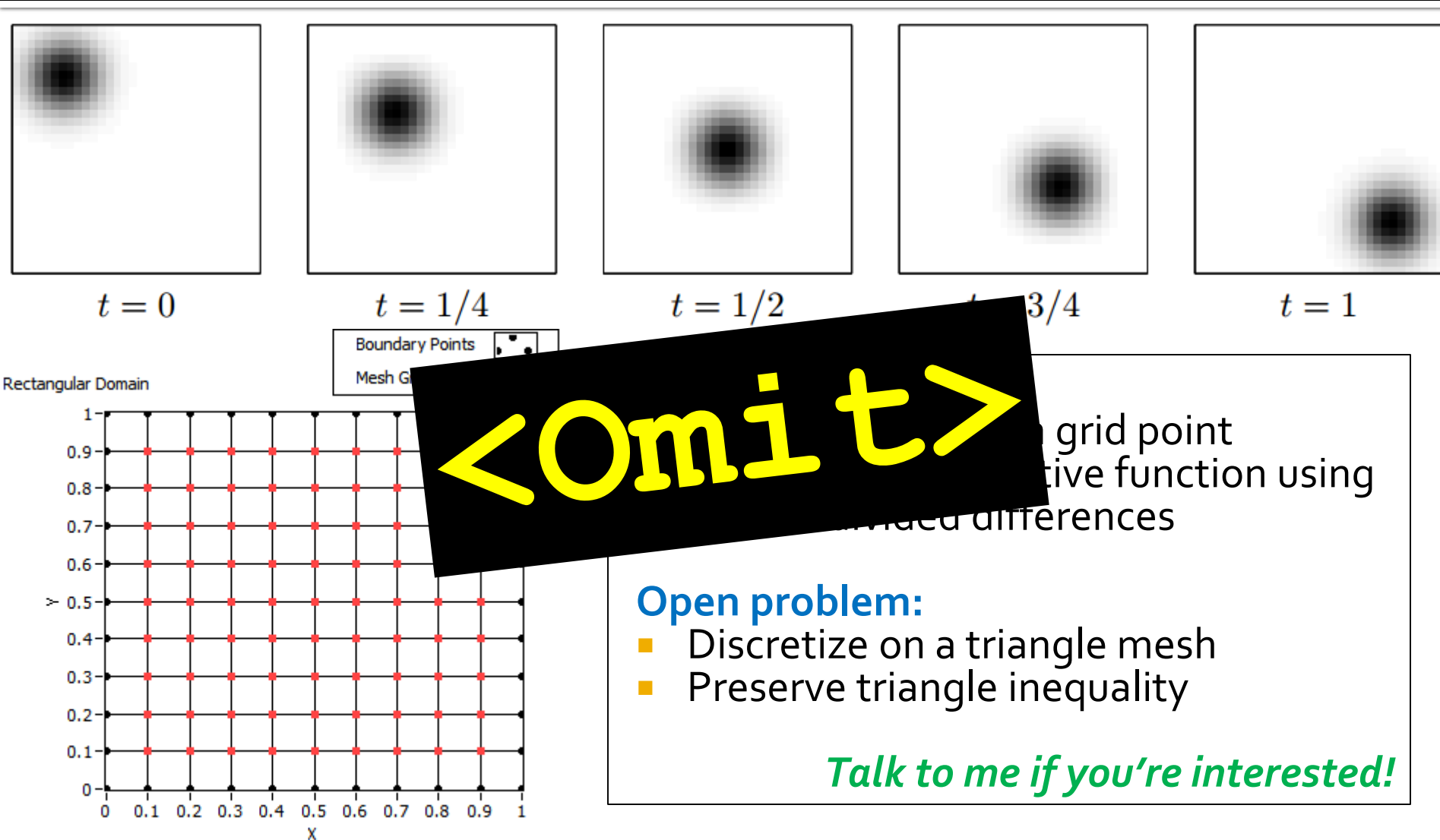
3. Discrete/discretized transport

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Discretization

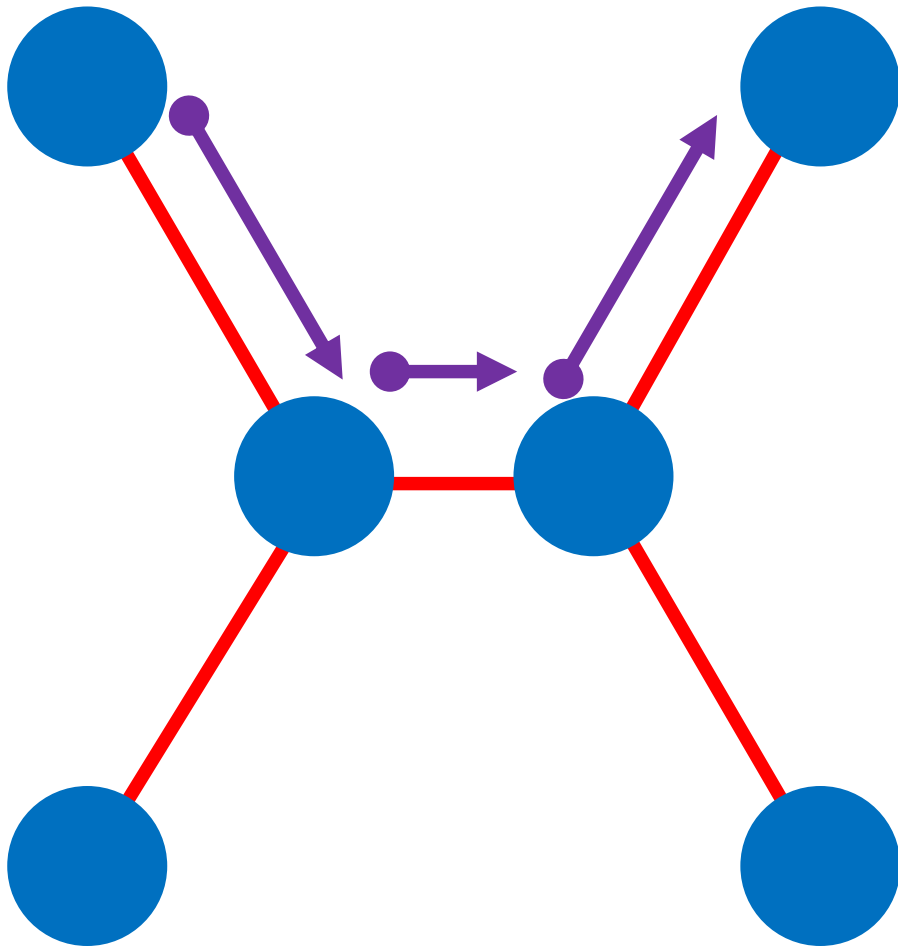


Open problem:

- Discretize on a triangle mesh
- Preserve triangle inequality

Talk to me if you're interested!

Beckmann Formulation



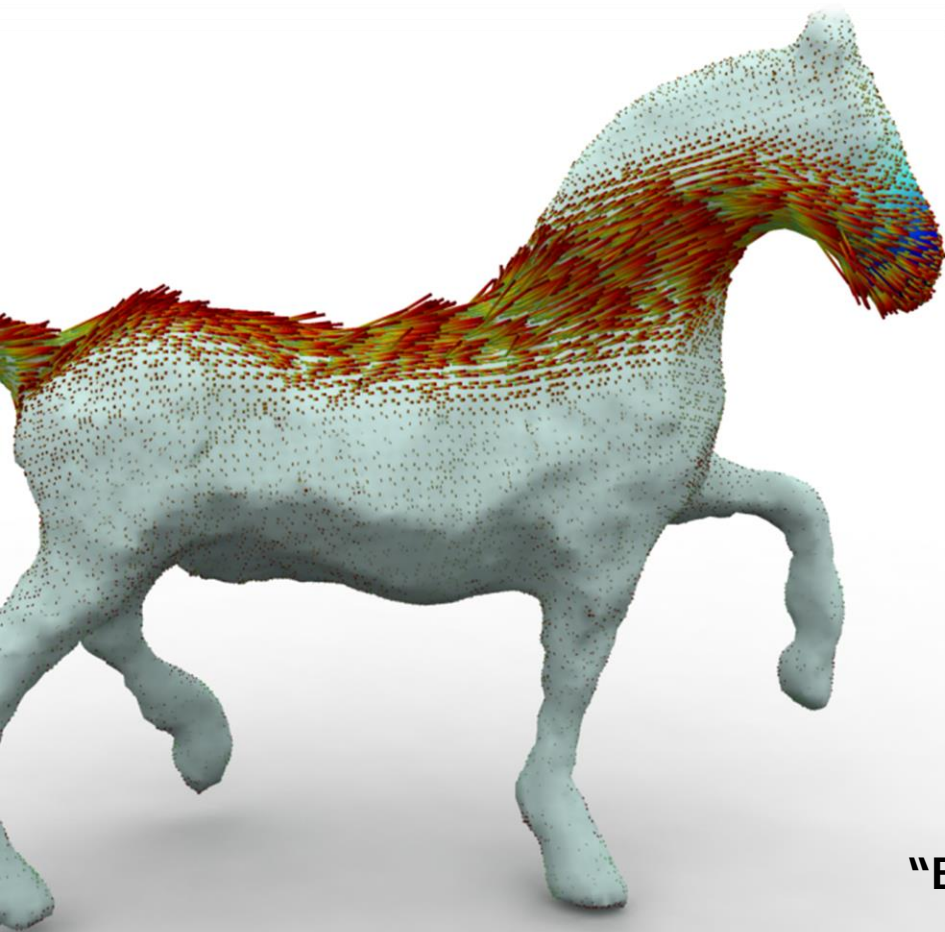
**Better scaling for
sparse graphs!**

$$\begin{array}{ll} \min_T & \sum_e c_e |J_e| \\ \text{s.t.} & D^\top J = \underbrace{p_1 - p_0}_f \end{array}$$

In computer science:

Network flow problem

Continuous Analog: Beckmann



Probabilities *advect*
along the surface

“Eulerian”

$$\mathcal{W}_1(\rho_0, \rho_1) = \begin{cases} \inf_J \int_M \|J(x)\| dx \\ \text{s.t. } \nabla \cdot J(x) = \rho_1(x) - \rho_0(x) \\ J(x) \cdot n(x) = 0 \quad \forall x \in \partial M \end{cases}$$

Solomon, Rustamov, Guibas, and Butscher.
“Earth Mover’s Distances on Discrete Surfaces.”

SIGGRAPH 2014

Plan For Today

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- Many formulas

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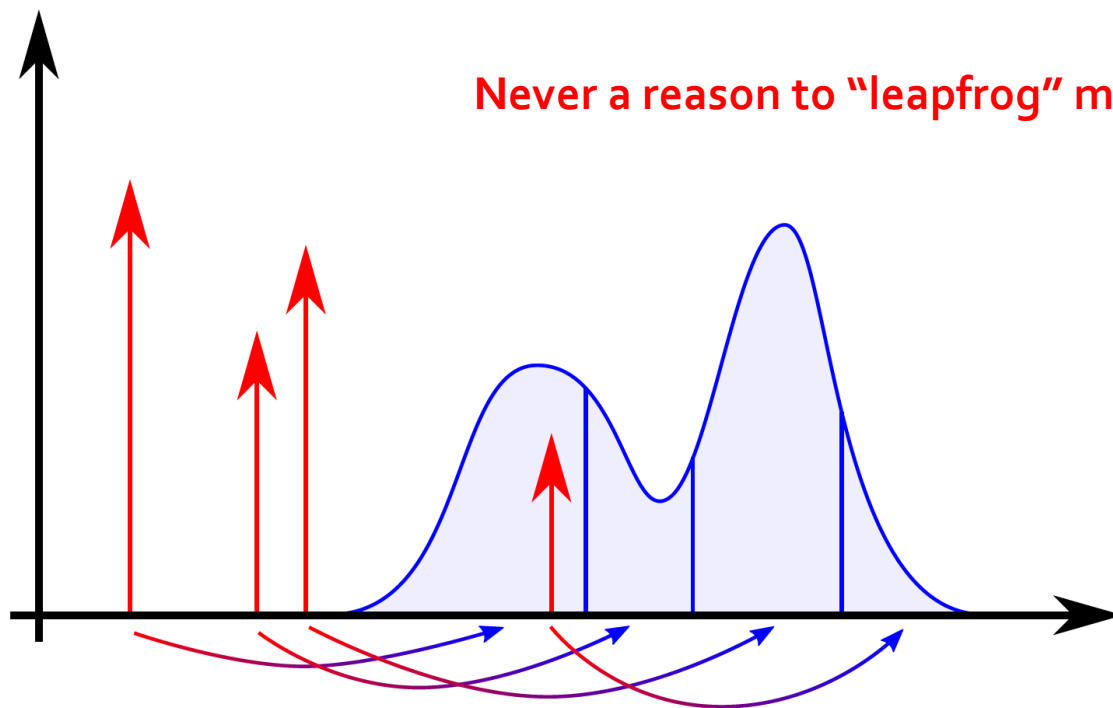


Recall:

Semidiscrete Transport

$$\mu_0 := \sum_{i=1}^{k_0} a_{0i} \delta_{x_{0i}}$$

$$\mu_1(S) := \int_S \rho_1(x) dx$$



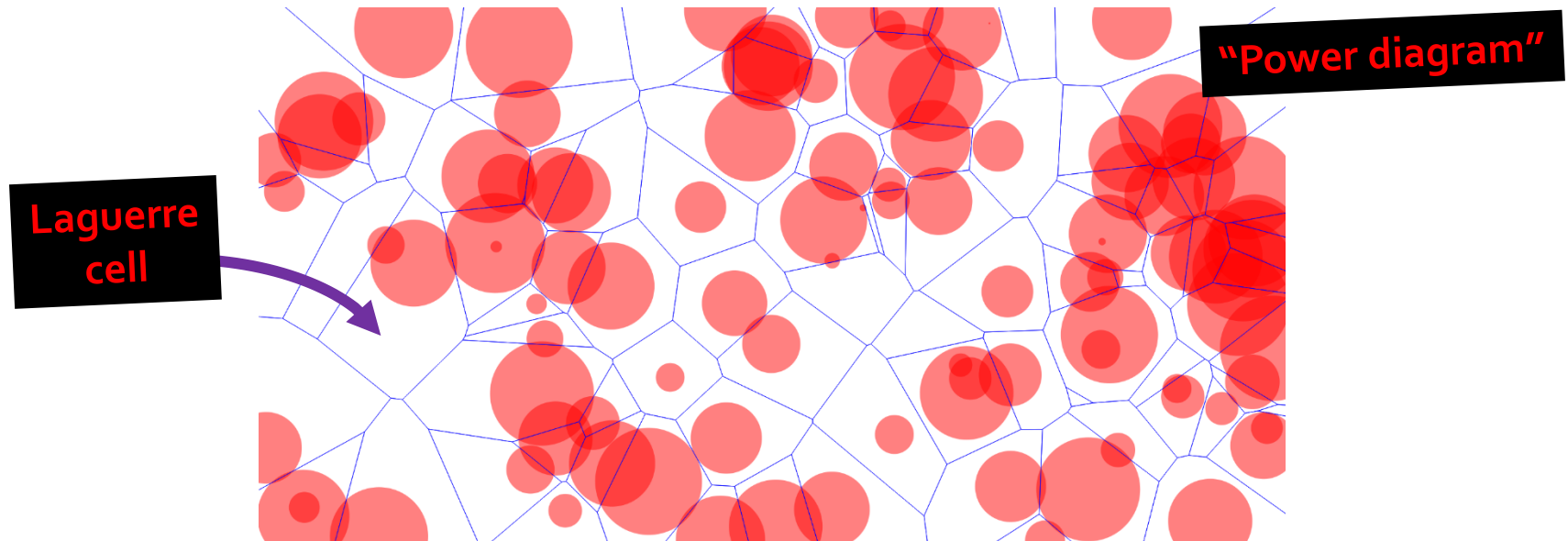
General Case: Semidiscrete

$$\mu_0 := \sum_{i=1}^k a_i \delta_{x_i}$$

$$\nu(S) := \int_S \rho(x) dx$$

$$\mathcal{W}_2^2(\mu, \nu) = \sup_{\phi \in \mathbb{R}^k} \sum_i \left[a_i \phi_i + \int_{\text{Lag}_\phi^c(x_i)} \rho(y) [c(x_i, y) - \phi_i] dA(y) \right]$$

$$\text{Lag}_\phi^c(x_i) := \{y \in \mathbb{R}^n : c(x_i, y) - \phi_i \leq c(x_j, y) - \phi_j \ \forall j \neq i\}$$



Semidiscrete Algorithm

$$F(\phi) := \sum_i \left[a_i \phi_i + \int_{\text{Lag}_\phi^c(x_i)} \rho(y) [c(x_i, y) - \phi_i] dA(y) \right]$$
$$\frac{\partial F}{\partial \phi_i} = a_i - \int_{\text{Lag}_\phi^c(x_i)} \rho(y) dA(y)$$

Concave in ϕ !

- **Simple algorithm:** Gradient ascent
Ingredients: Power diagram
- **More complex:** Newton's method
Converges globally [Kitagawa, Mériqot, & Thibert 2016]

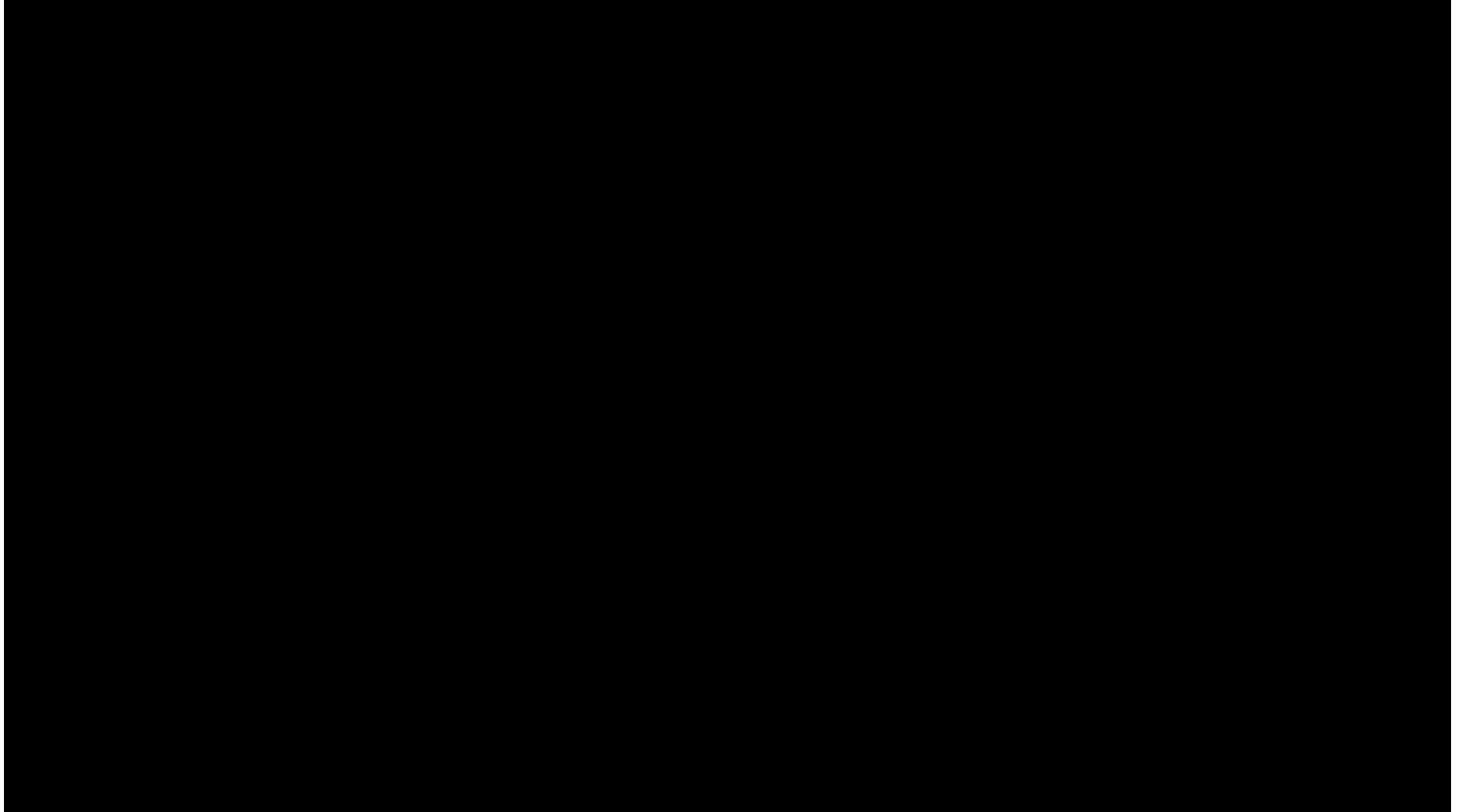
Application



Lévy. "A numerical algorithm for L2 semi-discrete optimal transport in 3D." (2014)

Points to tetrahedra

Application



Redux

Method	Advantages	Disadvantages
Entropic regularization	<ul style="list-style-type: none">• Fast• Easy to implement• Works on mesh using heat kernel	<ul style="list-style-type: none">• Blurry• Becomes singular as $\alpha \rightarrow 0$
Eulerian optimization	<ul style="list-style-type: none">• Provides displacement interpolation• Connection to PDE	<ul style="list-style-type: none">• Hard to optimize• Triangle mesh formulation unclear
Semidiscrete optimization	<ul style="list-style-type: none">• No regularization• Connection to “classical” geometry	<ul style="list-style-type: none">• Expensive computational geometry algorithms

Many others:
Stochastic transport, dual ascent, Monge-Ampère PDE, ...

Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

2. Applications

3. Discrete/discretized transport

- Entropic regularization
- Eulerian transport
- Semidiscrete transport

4. Extensions & frontiers



Application 1:

Example: Averaging

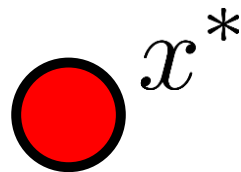
$$\text{Euclidean: } x^* := \left[\arg \min_{x \in \mathbb{R}^n} \sum_i \|x - x_i\|_2^2 \right] = \frac{1}{k} \sum_i x_i$$



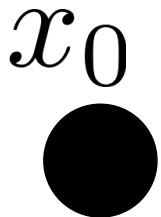
x_3



x_1



x^*

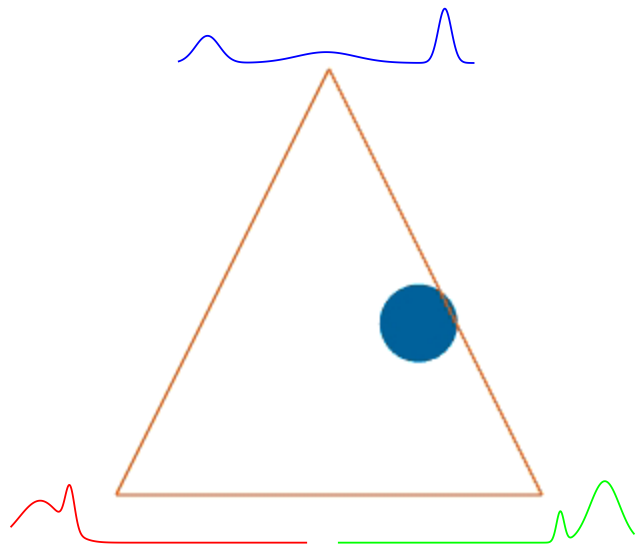


x_0

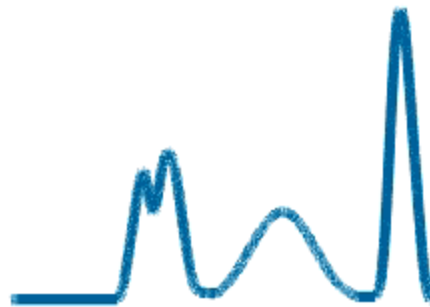


x_2

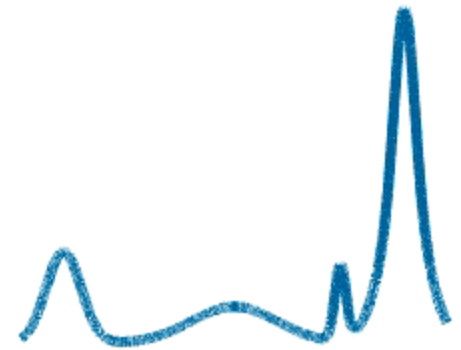
Barycenter Example



Wasserstein

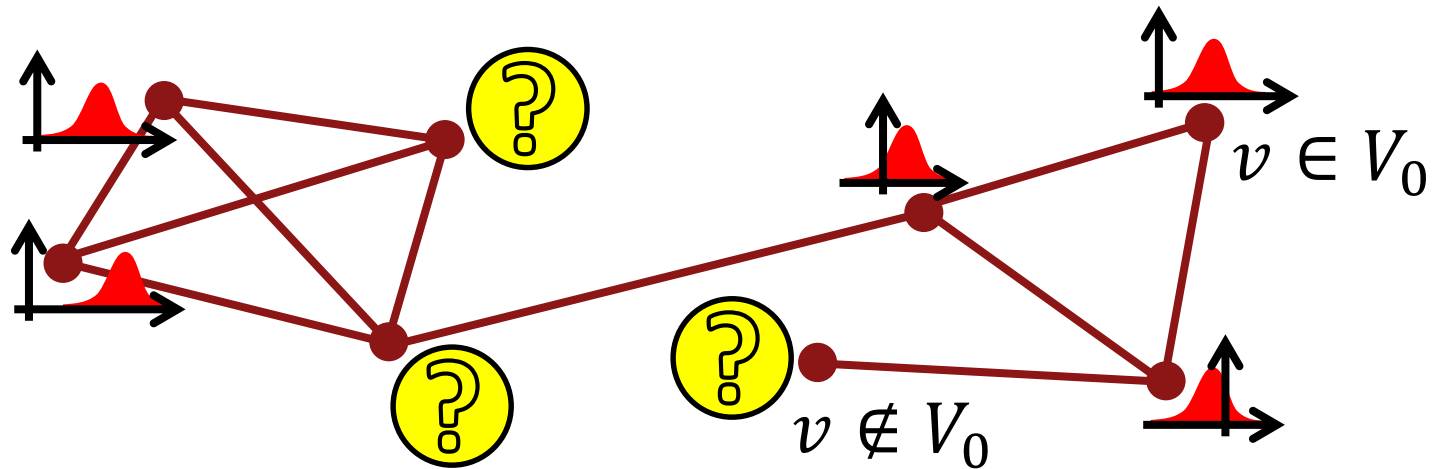


Euclidean

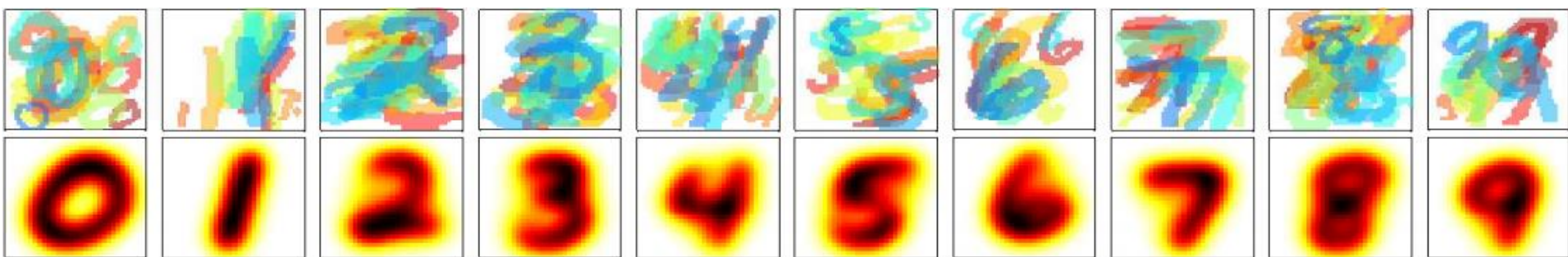


$$\text{Wasserstein: } \mu^* := \left[\arg \min_{\mu \in \text{Prob}(\mathbb{R}^n)} \sum_i \mathcal{W}_2^2(\mu, \mu_i) \right]$$

Barycenters in Machine Learning

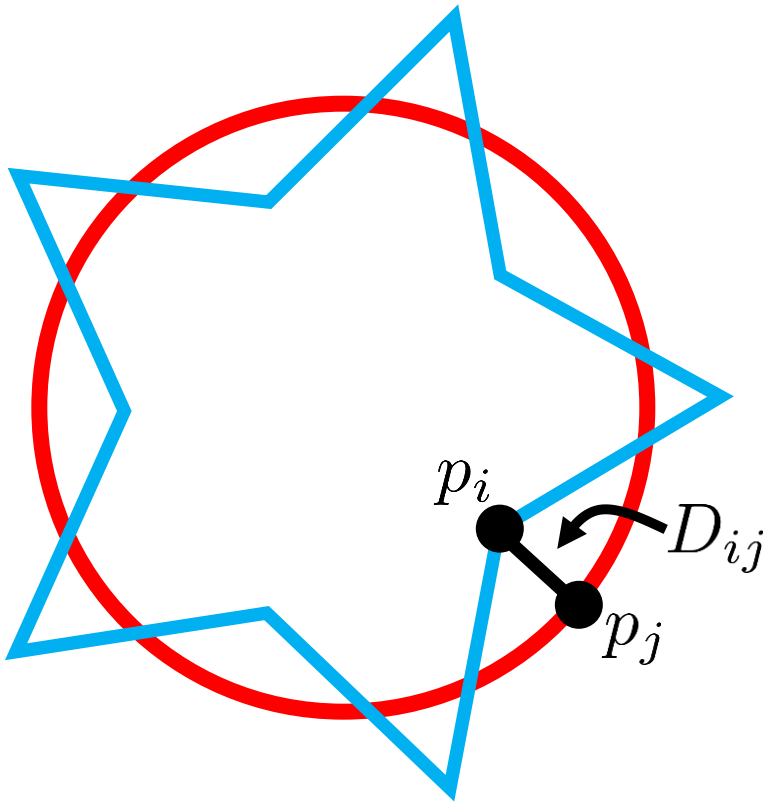


“Wasserstein Propagation for Semi-Supervised Learning” (Solomon et al.)



“Fast Computation of Wasserstein Barycenters” (Cuturi and Doucet)

Application 2: Model Problem: Linear Assignment



Between signals

$$\begin{aligned} \min_T \quad & \langle T, D \rangle \\ \text{s.t.} \quad & T \geq 0 \\ & T \mathbf{1} = \mathbf{1} \\ & T^\top \mathbf{1} = \mathbf{1} \end{aligned}$$

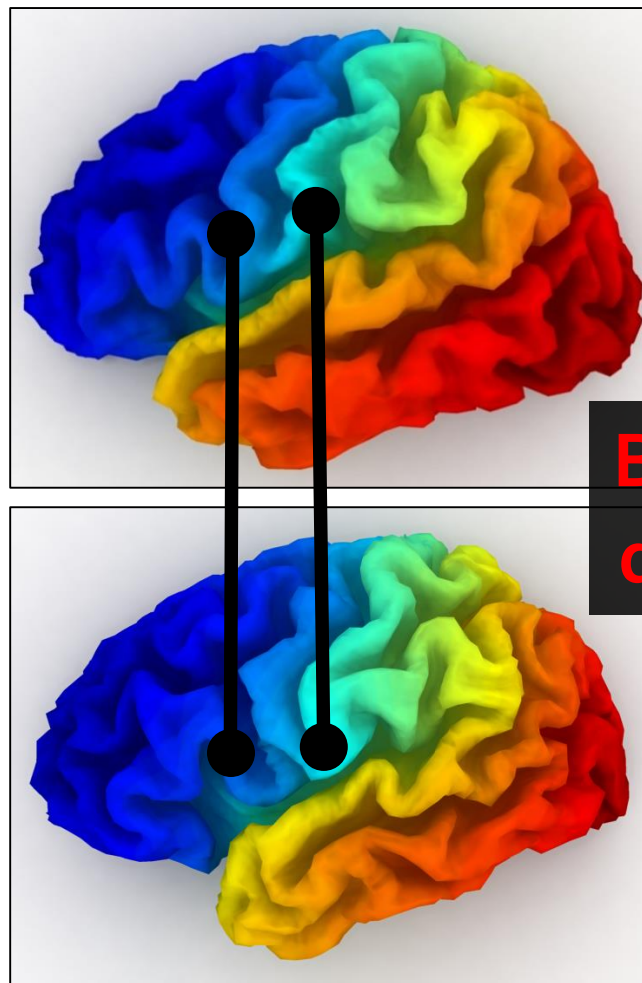
“No matched
point should
travel too far.”

Model Problem: Quadratic Matching

$$\begin{array}{ll}\min_T & \langle M_0 T, T M_1 \rangle \\ \text{s.t.} & T \geq 0 \\ & T \mathbf{1} = \mathbf{1} \\ & T^\top \mathbf{1} = \mathbf{1}\end{array}$$

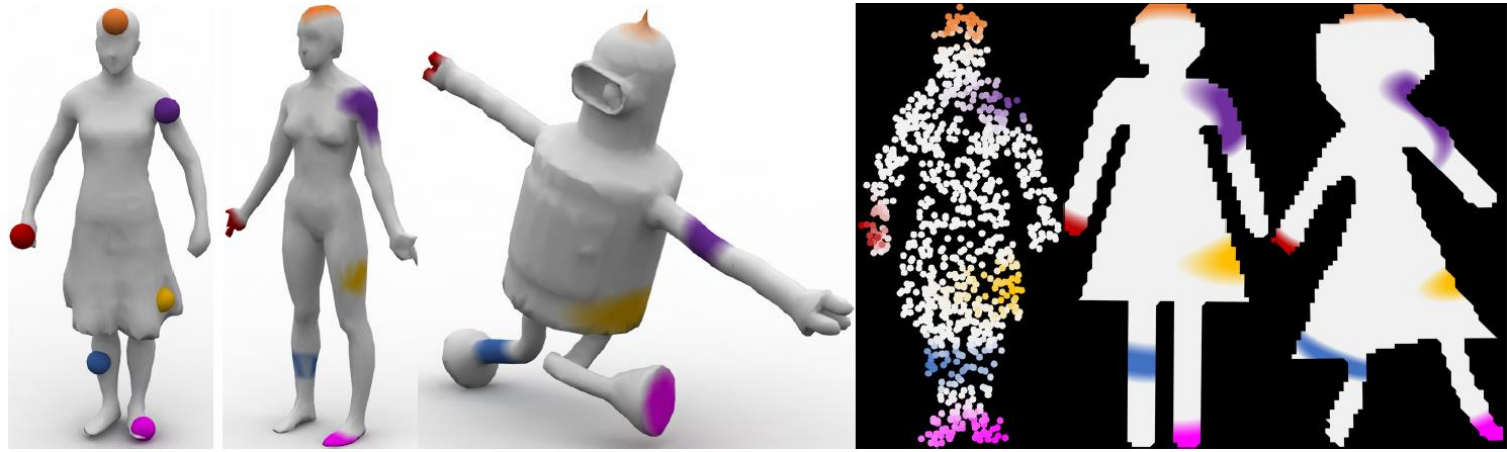
Nonconvex quadratic program!

“Nearby
points stay
nearby.”



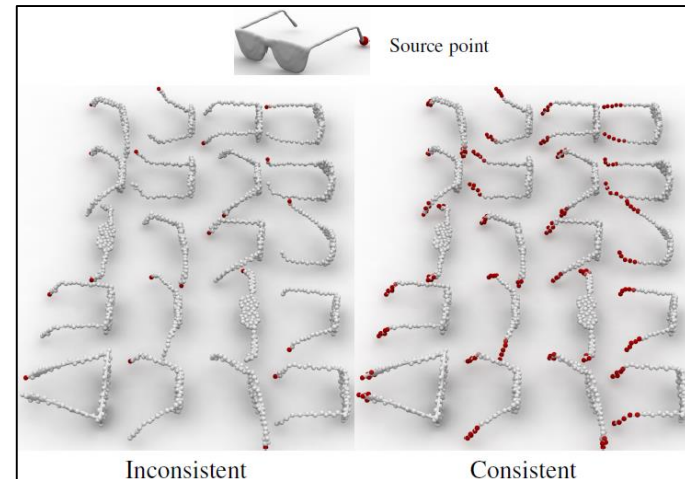
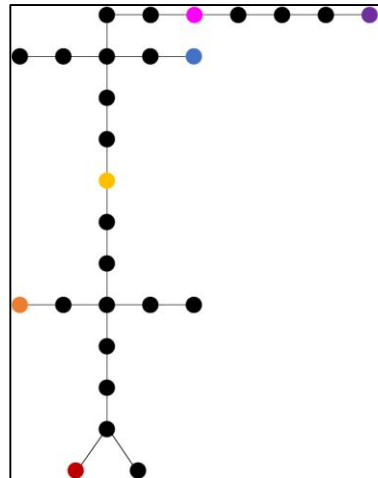
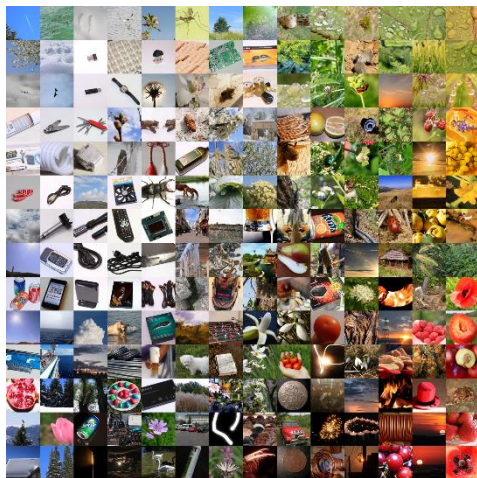
Between
domains

Variety of Correspondence Tasks



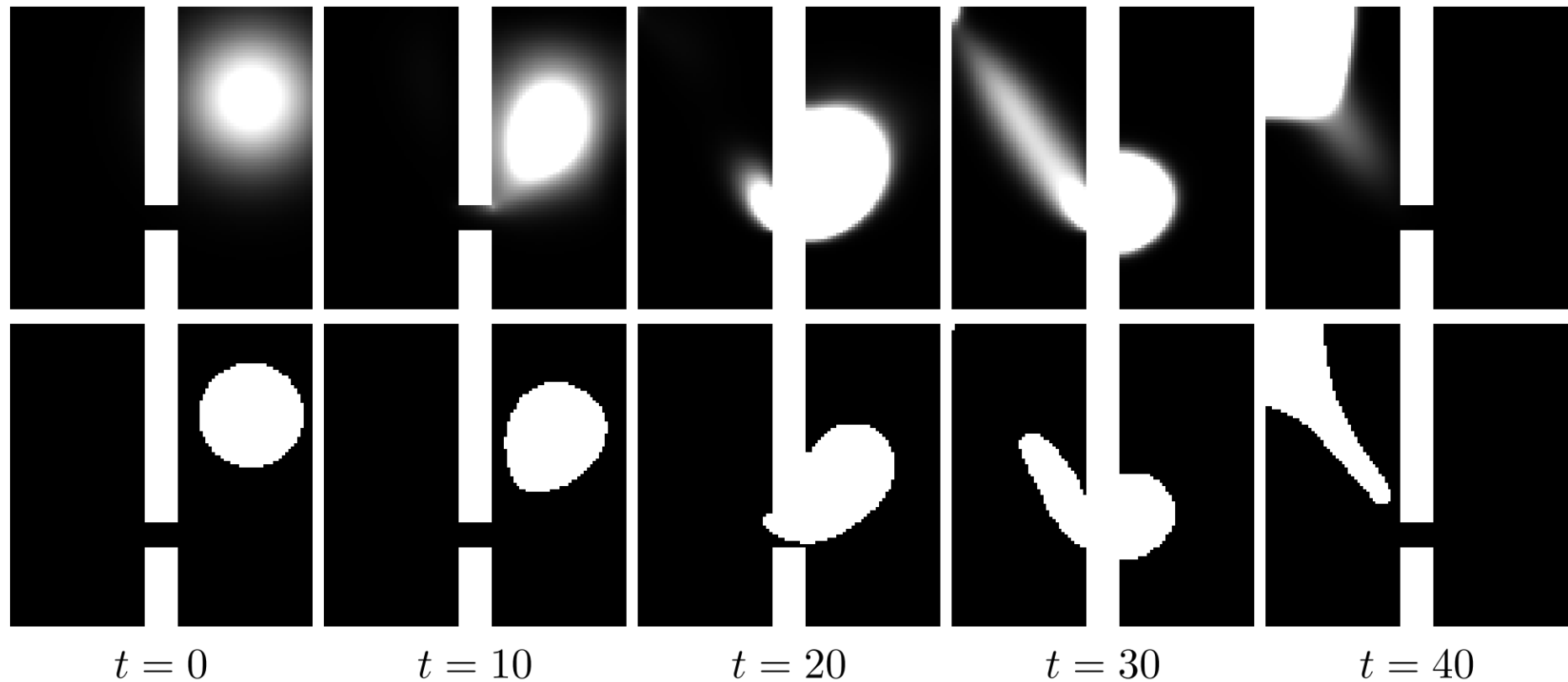
Source

Targets



Application 3:

Gradient Flows



“Entropic Wasserstein Gradient Flows” [Peyré 2015]

Interesting possibility for preserved structure!

Frontier:

Matrix Fields and Vector Measures

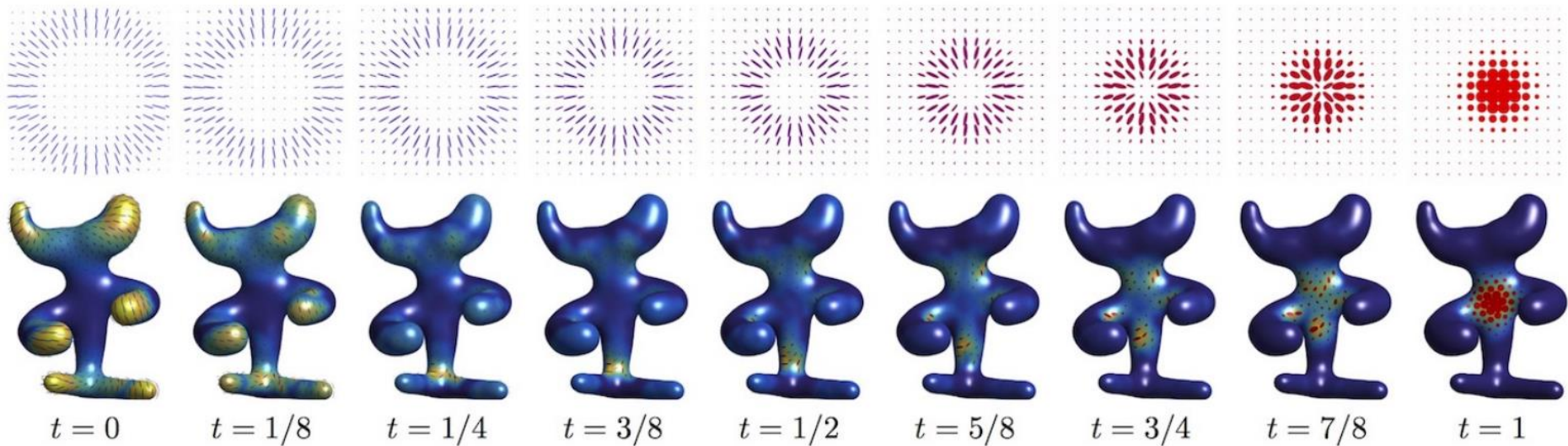


Image from

“Quantum Optimal Transport for Tensor Field Processing”

[Peyré et al. 2017]

Open problem: Dynamical version? Curved surfaces?

Frontier: Sampling Problems

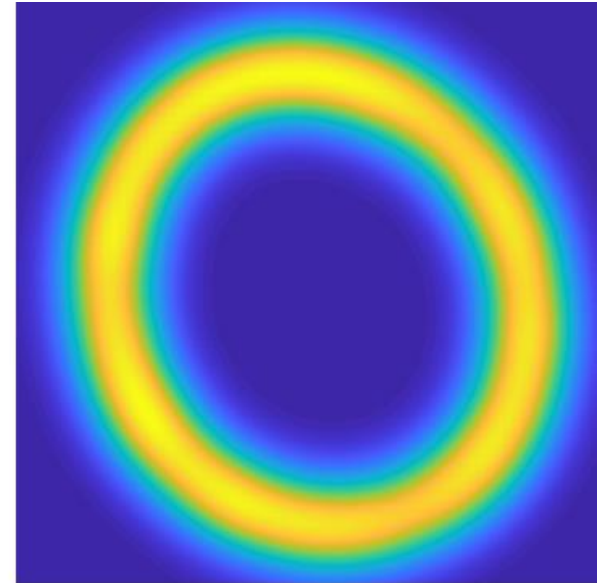
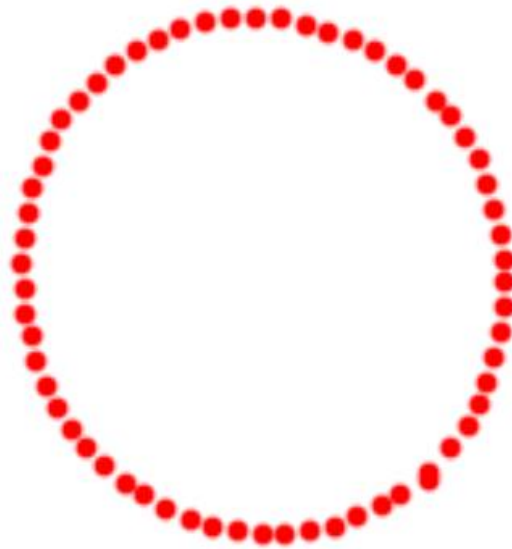
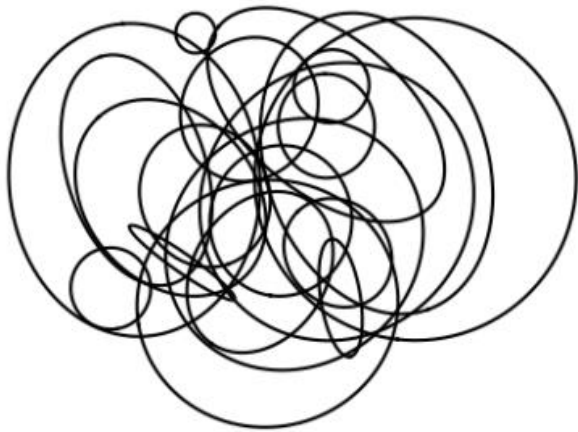
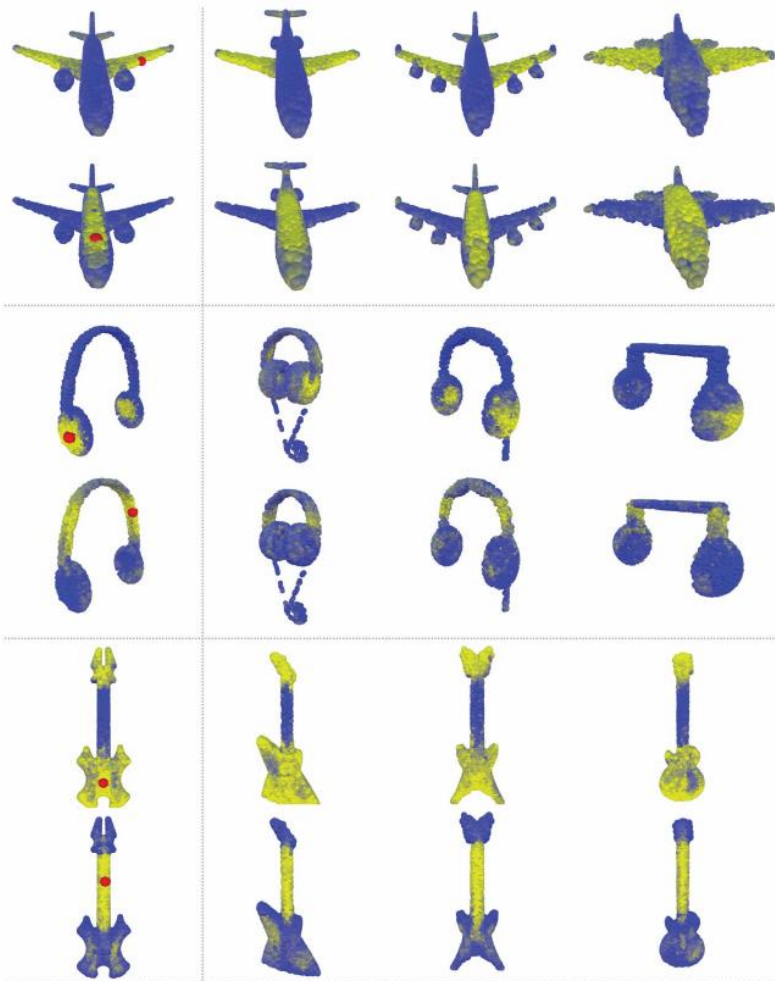


Image from
"Stochastic Wasserstein Barycenters"
[Claici et al. 2018]

Open problem: Sample from barycenter?

Frontier: Point Cloud Learning



Source points

Other point clouds from the same
category

Image from
"Dynamic Graph CNN for
Learning on Point Clouds"
[Wang et al. 2018]

**Open problem: Many repeated
instances of transport?**

Plan For Today

1. Introduction to optimal transport

- 1D examples
- Many formulas

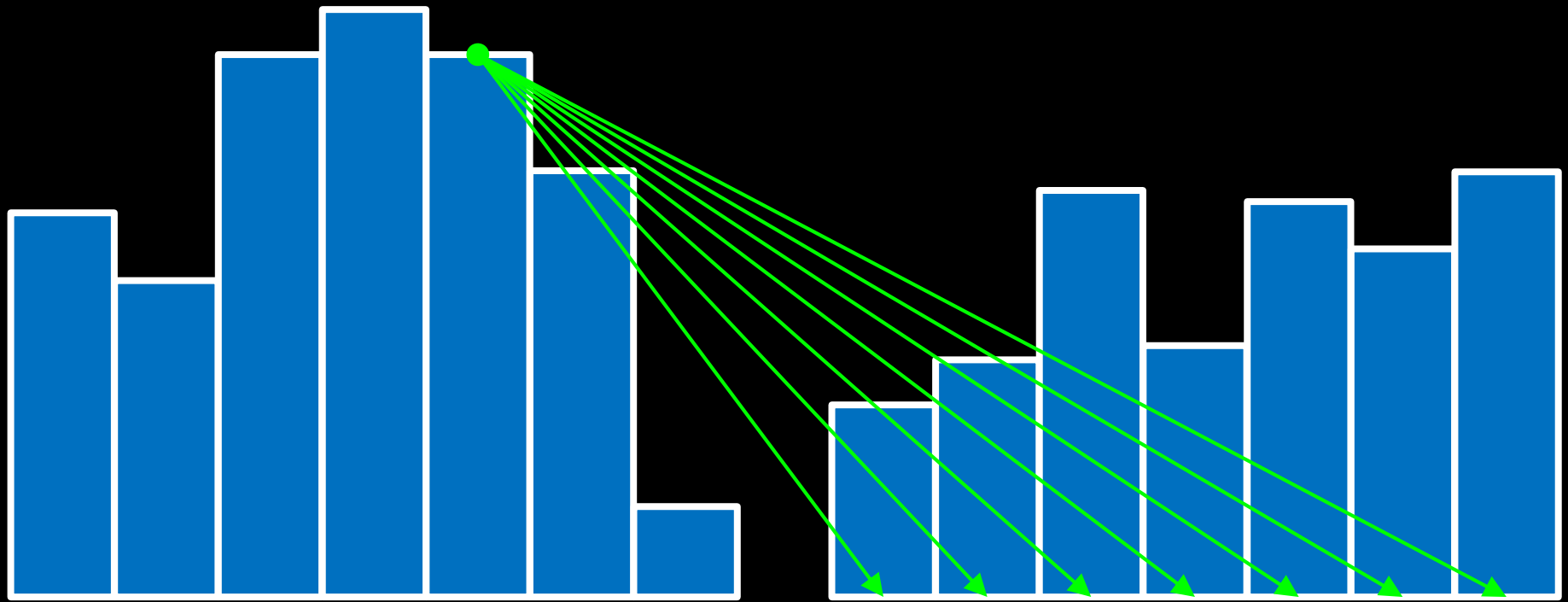
2. Applications

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Tutorial on
Optimal Transport

Questions?